

Analyzing Almost-Sure Termination of Probabilistic Term Rewriting via Innermost Almost-Sure Termination

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Research Group Computer Science 2
“Programming Languages and Verification”

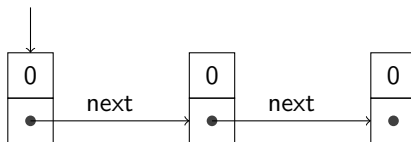
February 2023

Termination of TRSs

 $\mathcal{R}_{len}:$

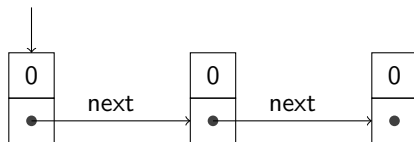
$$\begin{array}{ll} \text{len}(\text{nil}) & \rightarrow \mathcal{O} \\ \text{len}(\text{cons}(x, y)) & \rightarrow s(\text{len}(y)) \end{array}$$

Termination of TRSs

 $\mathcal{R}_{len}:$
 $len(nil) \rightarrow \mathcal{O}$
 $len(cons(x, y)) \rightarrow s(len(y))$


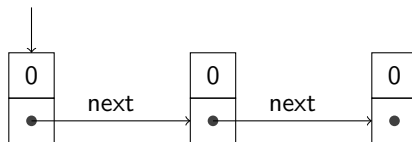
Termination of TRSs

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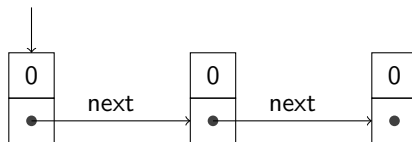
$$\text{len}(\text{cons}(\mathcal{O}, \text{cons}(\mathcal{O}, \text{cons}(\mathcal{O}, \text{nil}))))$$

Termination of TRSs

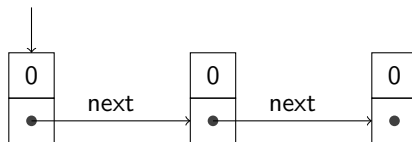
 $\mathcal{R}_{len}:$
 $len(nil) \rightarrow \mathcal{O}$
 $len(cons(x, y)) \rightarrow s(len(y))$


$$\rightarrow_{\mathcal{R}_{len}} \begin{array}{l} len(cons(\mathcal{O}, cons(\mathcal{O}, cons(\mathcal{O}, nil)))) \\ s(len(cons(\mathcal{O}, cons(\mathcal{O}, nil)))) \end{array}$$

Termination of TRSs

 $\mathcal{R}_{len}:$
 $len(nil) \rightarrow \mathcal{O}$
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 $len(cons(\mathcal{O}, cons(\mathcal{O}, cons(\mathcal{O}, nil))))$
 $\rightarrow_{\mathcal{R}_{len}}$
 $s(len(cons(\mathcal{O}, cons(\mathcal{O}, nil))))$
 $\rightarrow_{\mathcal{R}_{len}}$
 $s(s(len(cons(\mathcal{O}, nil))))$

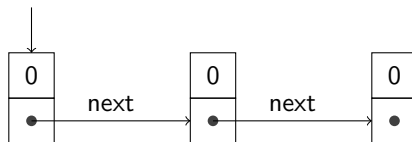
Termination of TRSs

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 $len(nil) \rightarrow \mathcal{O}$
 $len(cons(x, y)) \rightarrow s(len(y))$


$$\begin{aligned}
 & len(cons(\mathcal{O}, cons(\mathcal{O}, cons(\mathcal{O}, nil)))) \\
 \rightarrow_{\mathcal{R}_{len}} & s(len(cons(\mathcal{O}, cons(\mathcal{O}, nil)))) \\
 \rightarrow_{\mathcal{R}_{len}} & s(s(len(cons(\mathcal{O}, nil)))) \\
 \rightarrow_{\mathcal{R}_{len}} & s(s(s(len(nil))))
 \end{aligned}$$

Termination of TRSs

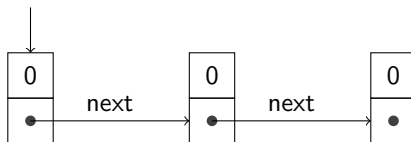
 $\mathcal{R}_{len}:$

$$\begin{aligned} \text{len}(\text{nil}) &\rightarrow \mathcal{O} \\ \text{len}(\text{cons}(x, y)) &\rightarrow s(\text{len}(y)) \end{aligned}$$


$$\begin{aligned} &\text{len}(\text{cons}(\mathcal{O}, \text{cons}(\mathcal{O}, \text{cons}(\mathcal{O}, \text{nil})))) \\ \rightarrow_{\mathcal{R}_{len}} &s(\text{len}(\text{cons}(\mathcal{O}, \text{cons}(\mathcal{O}, \text{nil})))) \\ \rightarrow_{\mathcal{R}_{len}} &s(s(\text{len}(\text{cons}(\mathcal{O}, \text{nil})))) \\ \rightarrow_{\mathcal{R}_{len}} &s(s(s(\text{len}(\text{nil})))) \\ \rightarrow_{\mathcal{R}_{len}} &s(s(s(\mathcal{O}))) \end{aligned}$$

Termination of TRSs

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$$\begin{aligned} &\text{len}(\text{cons}(\mathcal{O}, \text{cons}(\mathcal{O}, \text{cons}(\mathcal{O}, \text{nil})))) \\ \rightarrow_{\mathcal{R}_{len}} &s(\text{len}(\text{cons}(\mathcal{O}, \text{cons}(\mathcal{O}, \text{nil})))) \\ \rightarrow_{\mathcal{R}_{len}} &s(s(\text{len}(\text{cons}(\mathcal{O}, \text{nil})))) \\ \rightarrow_{\mathcal{R}_{len}} &s(s(s(\text{len}(\text{nil})))) \\ \rightarrow_{\mathcal{R}_{len}} &s(s(s(\mathcal{O}))) \end{aligned}$$

\mathcal{R} is terminating iff there is no infinite evaluation $t_0 \rightarrow_{\mathcal{R}} t_1 \rightarrow_{\mathcal{R}} \dots$

Where do we use TRSs?

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- Turing complete programming language

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- Turing complete programming language
- Backend language of tools for the analysis of programming languages

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Java

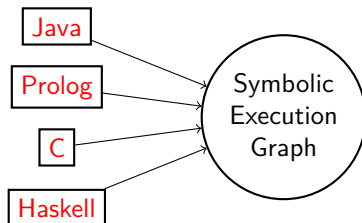
Prolog

C

Haskell

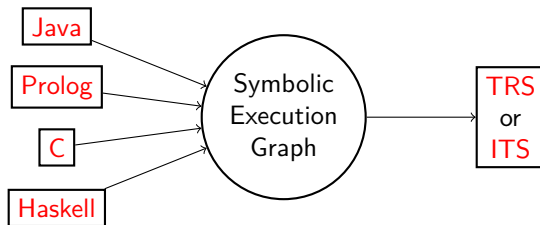
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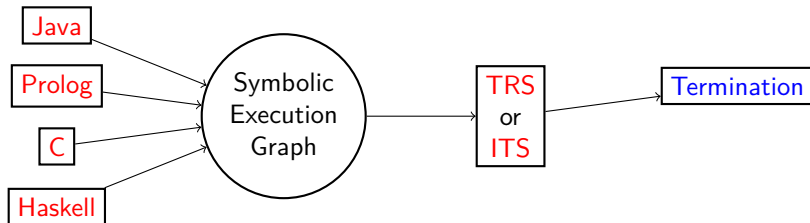
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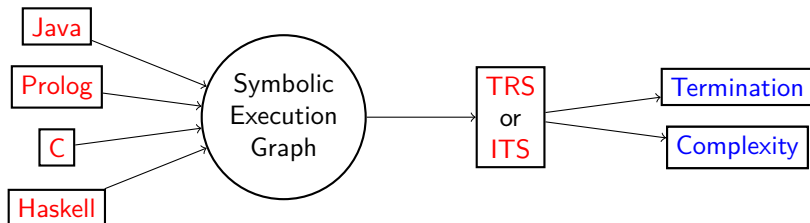
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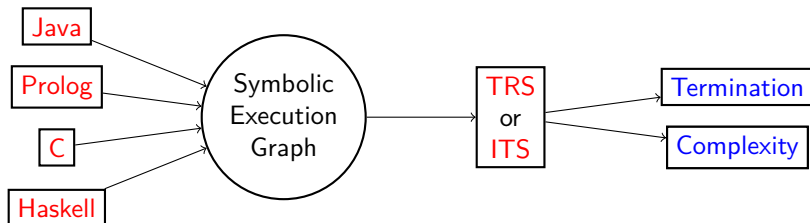
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- TRS: especially good for the analysis of algorithms concerning algebraic/user-defined data structures (lists, graphs, etc.)

Non-Determinism and Evaluation Strategies

\mathcal{R}_{plus} :

$$\begin{array}{lll} \text{plus}(\mathcal{O}, y) & \rightarrow & y \\ \text{plus}(s(x), y) & \rightarrow & s(\text{plus}(x, y)) \end{array}$$

Non-Determinism and Evaluation Strategies

\mathcal{R}_{plus} :

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$\text{plus}(s(\mathcal{O}), \text{plus}(\mathcal{O}, \mathcal{O}))$

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$\text{plus}(\text{s}(\mathcal{O}), \text{plus}(\mathcal{O}, \mathcal{O}))$



$\text{s}(\text{plus}(\mathcal{O}, \text{plus}(\mathcal{O}, \mathcal{O})))$

Non-Determinism and Evaluation Strategies

\mathcal{R}_{plus} :

$plus(\mathcal{O}, y) \rightarrow y$
 $plus(s(x), y) \rightarrow s(plus(x, y))$

$plus(s(\mathcal{O}), plus(\mathcal{O}, \mathcal{O}))$



$s(plus(\mathcal{O}, plus(\mathcal{O}, \mathcal{O})))$

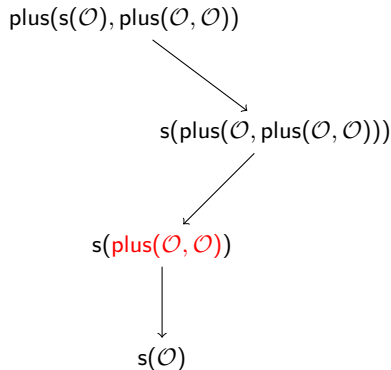


$s(plus(\mathcal{O}, \mathcal{O}))$

Non-Determinism and Evaluation Strategies

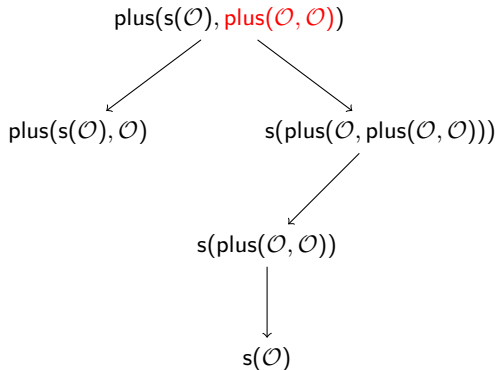
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Non-Determinism and Evaluation Strategies

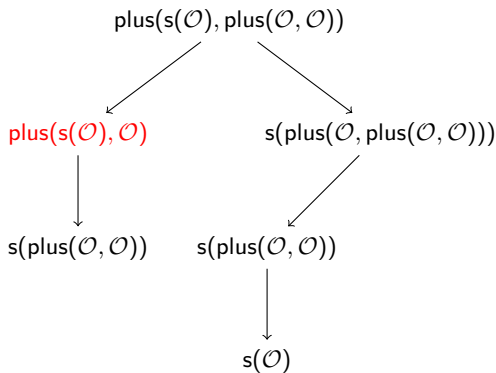
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Non-Determinism and Evaluation Strategies

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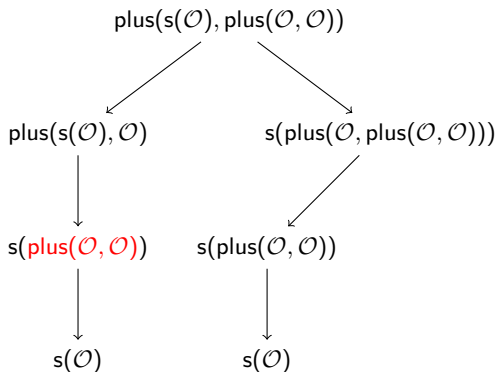
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Non-Determinism and Evaluation Strategies

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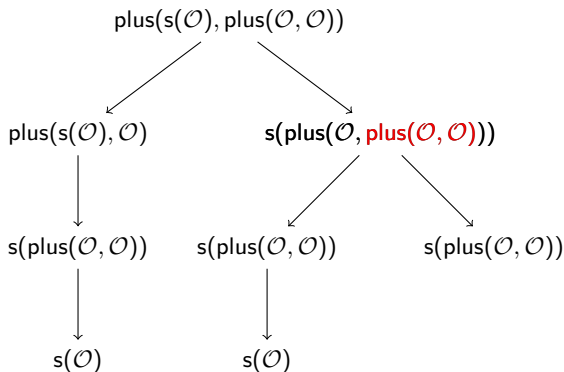
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Non-Determinism and Evaluation Strategies

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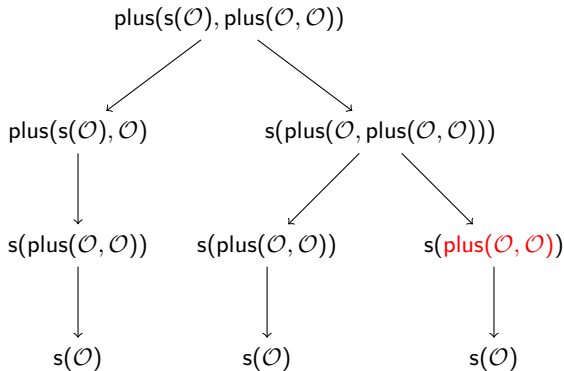
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Non-Determinism and Evaluation Strategies

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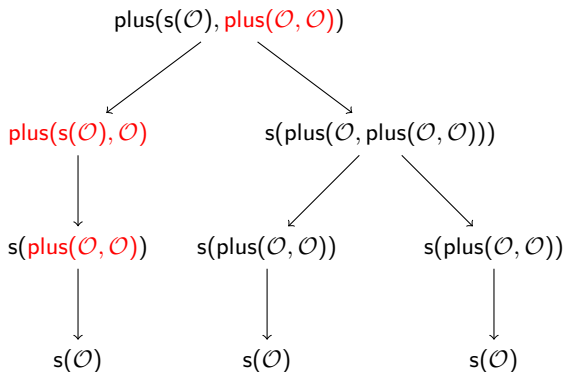
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Non-Determinism and Evaluation Strategies

 \mathcal{R}_{plus} :

$$\begin{aligned} \text{plus}(\mathcal{O}, y) &\rightarrow y \\ \text{plus}(s(x), y) &\rightarrow s(\text{plus}(x, y)) \end{aligned}$$



Innermost evaluation: always use an innermost reducible expression

Innermost Termination vs. Termination

$$\begin{array}{lcl} \mathcal{R}_1: & f(a, b, x) & \rightarrow f(x, x, x) \\ & g(x, y) & \rightarrow x \\ & g(x, y) & \rightarrow y \end{array}$$

Innermost Termination vs. Termination

$$\begin{array}{lll} \mathcal{R}_1: & f(a, b, x) & \rightarrow f(x, x, x) \\ & g(x, y) & \rightarrow x \\ & g(x, y) & \rightarrow y \end{array}$$

Not Terminating:

Innermost Termination vs. Termination

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Not Terminating:

$$f(a, b, g(a, b))$$

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Not Terminating:

$$f(a, b, g(a, b)) \rightarrow_{\mathcal{R}_1} f(g(a, b), g(a, b), g(a, b))$$

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But Innermost Terminating:

Innermost Termination vs. Termination

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But Innermost Terminating:

$$f(a, b, g(a, b))$$

Innermost Termination vs. Termination

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But Innermost Terminating:

$$f(a, b, g(a, b)) \xrightarrow{i} \mathcal{R}_1 f(a, b, a)$$

Innermost Termination vs. Termination

$$\begin{array}{lcl} \mathcal{R}_1: & f(a, b, x) & \rightarrow f(x, x, x) \\ & g(x, y) & \rightarrow x \\ & g(x, y) & \rightarrow y \end{array}$$

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But Innermost Terminating:

$$f(a, b, g(a, b)) \xrightarrow{i}_{\mathcal{R}_1} f(a, b, a) \xrightarrow{i}_{\mathcal{R}_1} f(a, a, a)$$

Innermost Termination vs. Termination

$$\begin{array}{lll} \mathcal{R}_1: & f(a, b, x) & \rightarrow f(x, x, x) \\ & g(x, y) & \rightarrow x \\ & g(x, y) & \rightarrow y \end{array}$$

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But Innermost Terminating:

$$f(a, b, g(a, b)) \xrightarrow{i}_{\mathcal{R}_1} f(a, b, a) \xrightarrow{i}_{\mathcal{R}_1} f(a, a, a) \leftarrow \text{normal form}$$

Innermost Termination vs. Termination

$$\begin{array}{lll} \mathcal{R}_1: & f(a, b, x) & \rightarrow f(x, x, x) \\ & g(x, y) & \rightarrow x \\ & g(x, y) & \rightarrow y \end{array}$$

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But Innermost Terminating:

$$f(a, b, g(a, b)) \xrightarrow{i}_{\mathcal{R}_1} f(a, b, a) \xrightarrow{i}_{\mathcal{R}_1} f(a, a, a) \leftarrow \text{normal form}$$

$$f(a, b, g(a, b))$$

Innermost Termination vs. Termination

$$\begin{array}{lll} \mathcal{R}_1: & f(a, b, x) & \rightarrow f(x, x, x) \\ & g(x, y) & \rightarrow x \\ & g(x, y) & \rightarrow y \end{array}$$

Not Terminating:

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But Innermost Terminating:

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$$f(a, b, g(a, b)) \xrightarrow{i}_{\mathcal{R}_1} f(a, b, b)$$

Innermost Termination vs. Termination

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$$f(a, b, g(a, b)) \xrightarrow{i}_{\mathcal{R}_1} f(a, b, b) \xrightarrow{i}_{\mathcal{R}_1} f(b, b, b)$$

Innermost Termination vs. Termination

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But Innermost Terminating:

$$f(a, b, g(a, b)) \xrightarrow{i}_{\mathcal{R}_1} f(a, b, a) \xrightarrow{i}_{\mathcal{R}_1} f(a, a, a) \leftarrow \text{normal form}$$

$$f(a, b, g(a, b)) \xrightarrow{i}_{\mathcal{R}_1} f(a, b, b) \xrightarrow{i}_{\mathcal{R}_1} f(b, b, b) \leftarrow \text{normal form}$$

Innermost Termination vs. (full) Termination

 \mathcal{R}_2 :

$$\begin{array}{lcl} f(a) & \rightarrow & f(a) \\ a & \rightarrow & b \end{array}$$

Innermost Termination vs. (full) Termination

 \mathcal{R}_2 :

$$\begin{array}{lcl} f(a) & \rightarrow & f(a) \\ a & \rightarrow & b \end{array}$$

Not Terminating:

Innermost Termination vs. (full) Termination

$$\mathcal{R}_2: \quad \begin{array}{lcl} f(a) & \rightarrow & f(a) \\ a & \rightarrow & b \end{array}$$

Not Terminating:

$f(a)$

Innermost Termination vs. (full) Termination

\mathcal{R}_2 :

$$\begin{array}{ccc} f(a) & \rightarrow & f(a) \\ a & \rightarrow & b \end{array}$$

Not Terminating:

$$f(a) \rightarrow_{\mathcal{R}_2} f(a)$$

Innermost Termination vs. (full) Termination

\mathcal{R}_2 :

$$\begin{array}{ccc} f(a) & \rightarrow & f(a) \\ a & \rightarrow & b \end{array}$$

Not Terminating:

$$f(a) \rightarrow_{\mathcal{R}_2} f(a) \rightarrow_{\mathcal{R}_2} \dots$$

Innermost Termination vs. (full) Termination

$$\mathcal{R}_2: \quad \begin{array}{lcl} f(a) & \rightarrow & f(a) \\ a & \rightarrow & b \end{array}$$

Not Terminating:

$$f(a) \rightarrow_{\mathcal{R}_2} f(a) \rightarrow_{\mathcal{R}_2} \dots$$

But Innermost Terminating:

Innermost Termination vs. (full) Termination

$$\mathcal{R}_2: \quad \begin{array}{ccc} f(a) & \rightarrow & f(a) \\ a & \rightarrow & b \end{array}$$

Not Terminating:

$$f(a) \rightarrow_{\mathcal{R}_2} f(a) \rightarrow_{\mathcal{R}_2} \dots$$

But Innermost Terminating:

$$f(a)$$

Innermost Termination vs. (full) Termination

$$\mathcal{R}_2: \quad \begin{array}{ccc} f(a) & \rightarrow & f(a) \\ a & \rightarrow & b \end{array}$$

Not Terminating:

$$f(a) \rightarrow_{\mathcal{R}_2} f(a) \rightarrow_{\mathcal{R}_2} \dots$$

But Innermost Terminating:

$$f(\textcolor{red}{a}) \xrightarrow{i}_{\mathcal{R}_2} f(b)$$

Innermost Termination vs. (full) Termination

$$\mathcal{R}_2: \quad \begin{array}{ccc} f(a) & \rightarrow & f(a) \\ a & \rightarrow & b \end{array}$$

Not Terminating:

$$f(a) \rightarrow_{\mathcal{R}_2} f(a) \rightarrow_{\mathcal{R}_2} \dots$$

But Innermost Terminating:

$$f(a) \xrightarrow{i}_{\mathcal{R}_2} f(b) \leftarrow \text{normal form}$$

Overlapping

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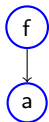
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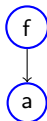


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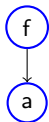
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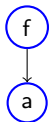


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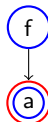
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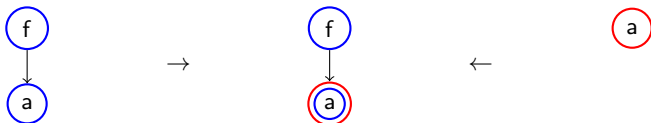


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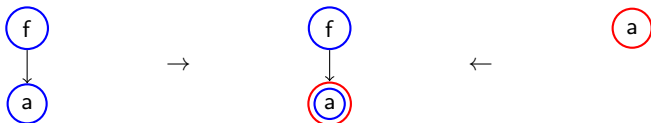
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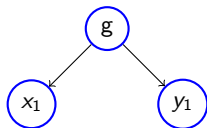
$\rightarrow \mathcal{R}_2$ is overlapping.

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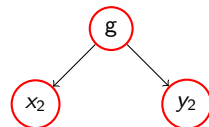
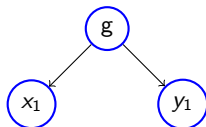
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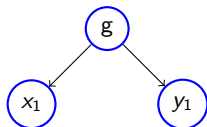
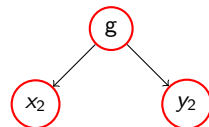
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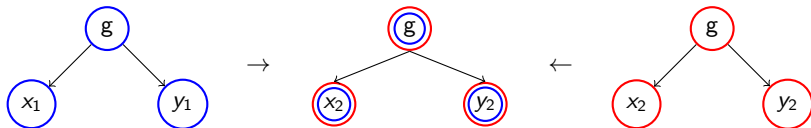
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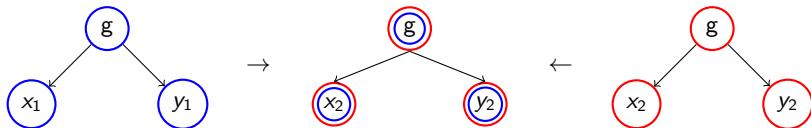
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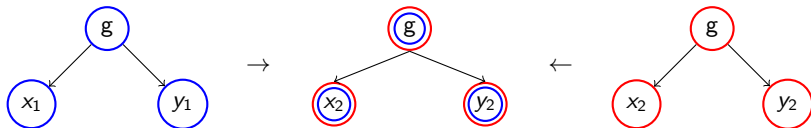
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$$\text{plus}(s(\mathcal{O}), \text{plus}(\mathcal{O}, \mathcal{O}))$$

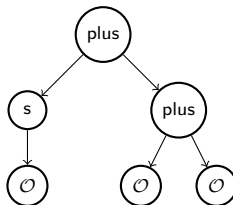
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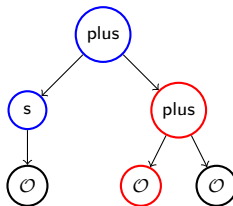
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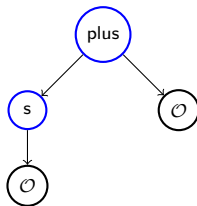
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Due to minimality rewriting $t_N \xrightarrow{\tau}_{\mathcal{R}} s_N$ at a position τ below π can not start an infinite evaluation

Proof cont.

From

$$t_N \rightarrow_{\mathcal{R}} t_{N+1} \rightarrow_{\mathcal{R}} t_{N+2} \rightarrow_{\mathcal{R}} \dots$$

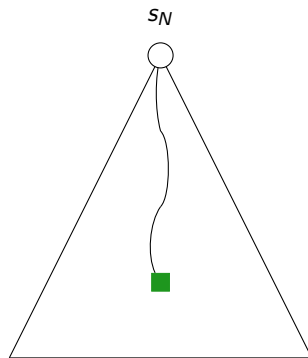
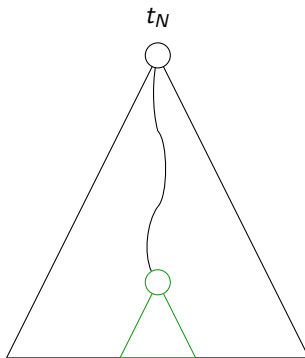
construct infinite evaluation

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Proof cont.

Rewriting possibilities:

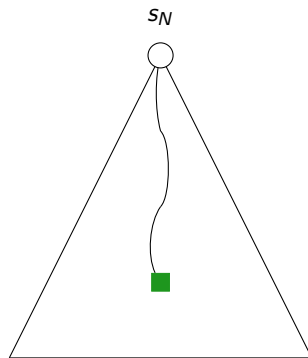
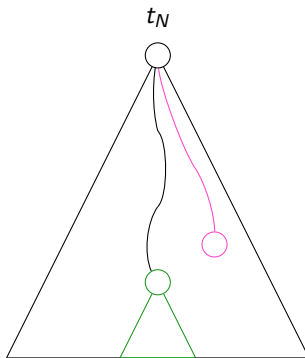
1. Rewriting at an orthogonal position
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Proof cont.

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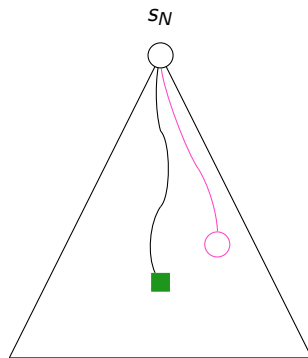
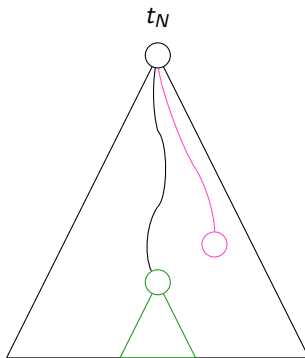
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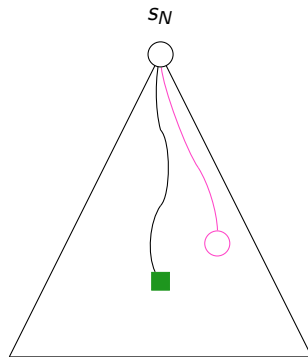
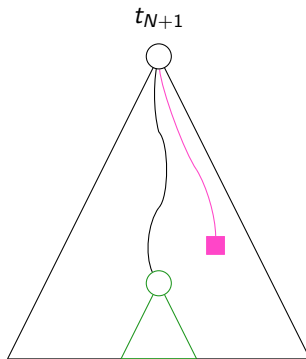
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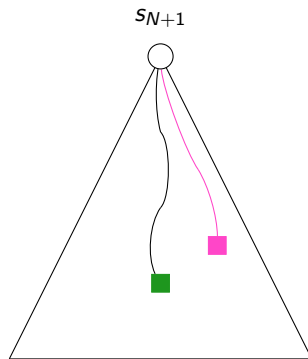
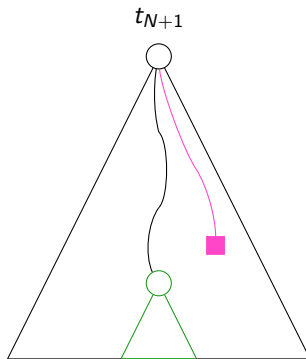
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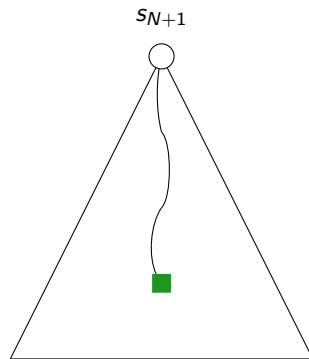
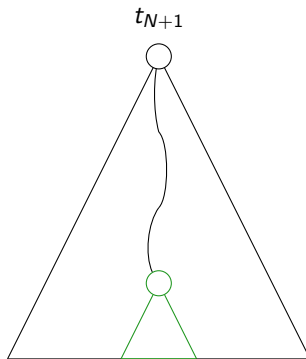
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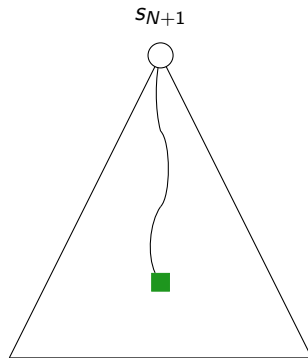
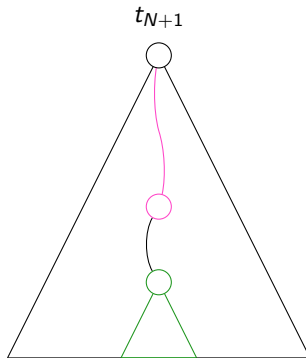
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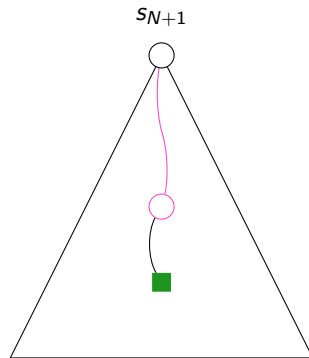
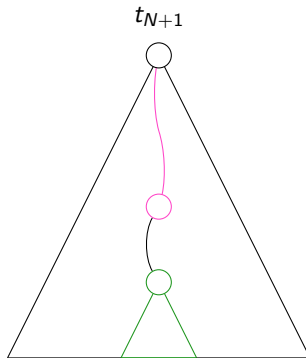
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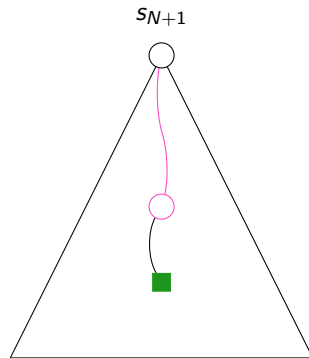
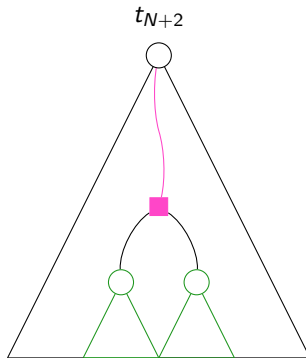
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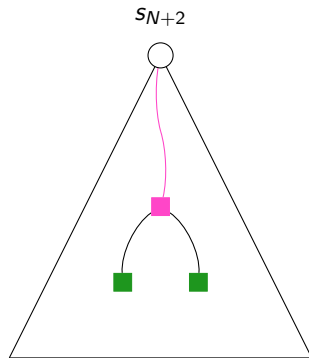
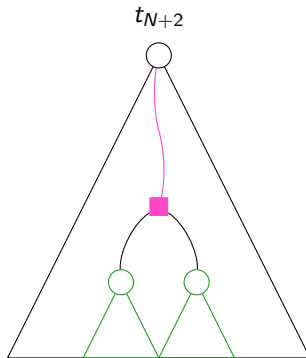
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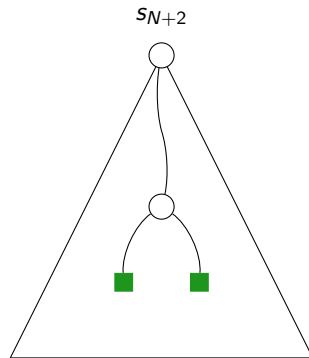
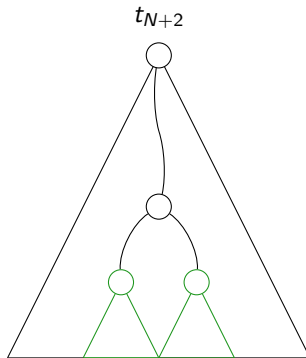
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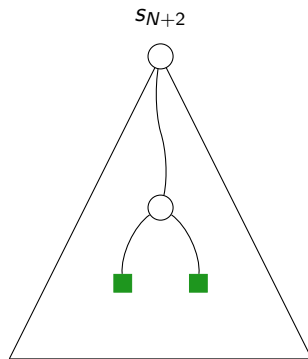
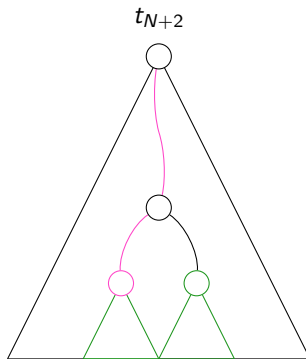
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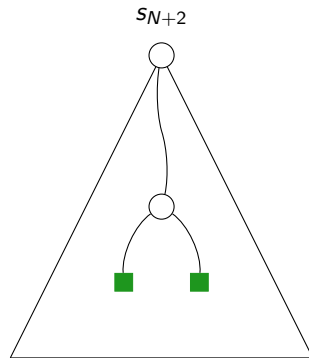
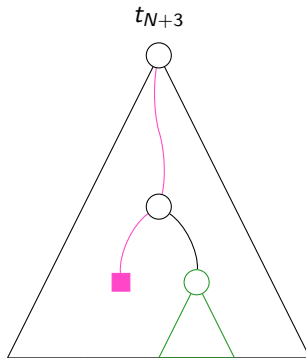
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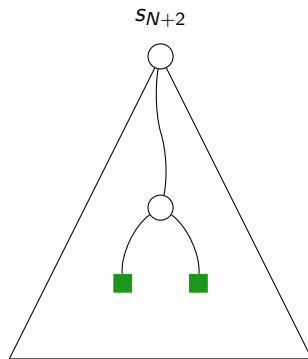
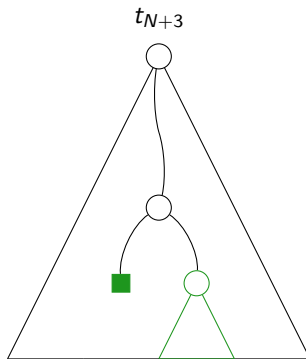
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Termination of PTRSs

$$\begin{array}{lll} \mathcal{S}_{len}: & \text{len}(\text{nil}) & \rightarrow \{1/2 : \mathcal{O}, 1/2 : \text{len}(\text{nil})\} \\ & \text{len}(\text{cons}(x, y)) & \rightarrow \{1/2 : s(\text{len}(y)), 1/2 : \text{len}(\text{cons}(x, y))\} \end{array}$$

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Computation of $\text{len}(\text{nil})$:

Termination of PTRSs

$$\begin{array}{ll} \mathcal{S}_{len}: & \text{len}(\text{nil}) \rightarrow \{1/2 : \mathcal{O}, 1/2 : \text{len}(\text{nil})\} \\ & \text{len}(\text{cons}(x, y)) \rightarrow \{1/2 : s(\text{len}(y)), 1/2 : \text{len}(\text{cons}(x, y))\} \end{array}$$

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- \mathcal{S}_{len} is AST

Innermost AST vs. (full) AST

\mathcal{S}_2 :

$f(a)$	\rightarrow	$\{1 : f(a)\}$
a	\rightarrow	$\{1 : b\}$

Innermost AST vs. (full) AST

\mathcal{S}_2 :

$$\begin{array}{lcl} f(a) & \rightarrow & \{1 : f(a)\} \\ a & \rightarrow & \{1 : b\} \end{array}$$

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Innermost AST vs. (full) AST

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$$\{1 : f(a)\} \Rightarrow_{\mathcal{S}_2} \{1 : f(a)\} \Rightarrow_{\mathcal{S}_2} \dots$$

But Innermost AST:

$$\{1 : f(a)\} \xRightarrow{i}_{\mathcal{S}_2} \{1 : f(b)\} \leftarrow \text{normal form}$$

Properties for Equality of Termination

Does non-overlapping still suffice?

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$$\mathcal{S}_3: \quad g \rightarrow \{3/4 : c(g, g), 1/4 : \perp\}$$

\rightarrow Biased random walk with $p = 3/4 > 1/2$, hence not AST.

Properties for Equality of Termination cont.

But **Innermost AST**:

$$\begin{aligned} g &\rightarrow \{^{3/4} : f(g), ^{1/4} : \perp\} \\ f(x) &\rightarrow \{1 : c(x, x)\} \end{aligned}$$

Properties for Equality of Termination cont.

But **Innermost AST**:

$$\begin{aligned} g &\rightarrow \{^{3/4} : f(g), ^{1/4} : \perp\} \\ f(x) &\rightarrow \{1 : c(x, x)\} \end{aligned}$$

$\mu_0 :$

$$1 : g$$

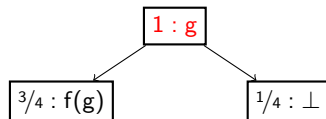
Properties for Equality of Termination cont.

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$\mu_0 :$

$\mu_1 :$



Properties for Equality of Termination cont.

But **Innermost AST**:

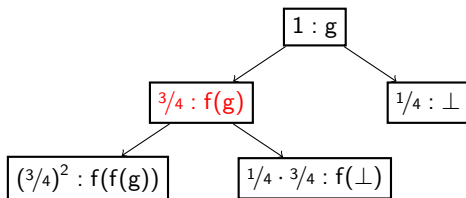
$$g \rightarrow \{3/4 : f(g), 1/4 : \perp\}$$

$$f(x) \rightarrow \{1 : c(x, x)\}$$

$\mu_0 :$

$\mu_1 :$

$\mu_2 :$

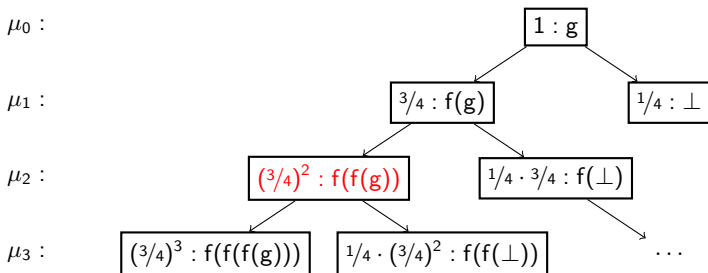


Properties for Equality of Termination cont.

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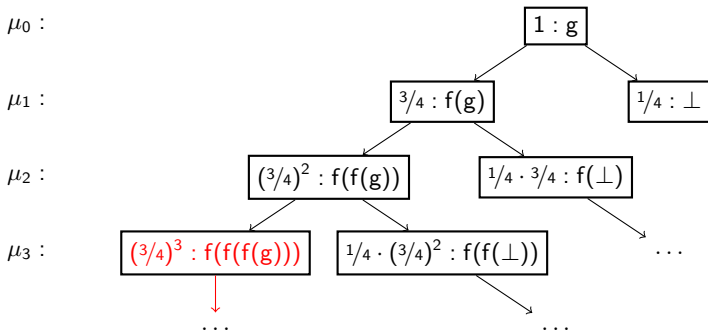


Properties for Equality of Termination cont.

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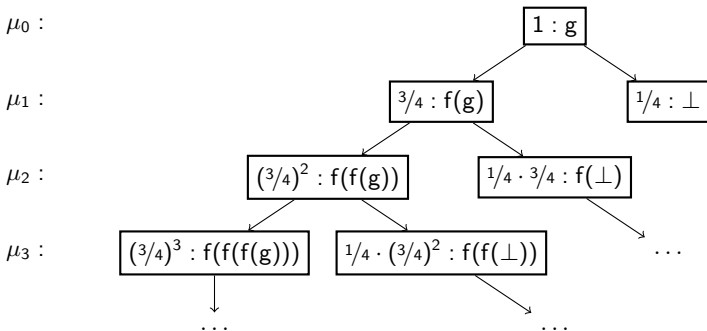
$$f(x) \rightarrow \{1 : c(x, x)\}$$



Properties for Equality of Termination cont.

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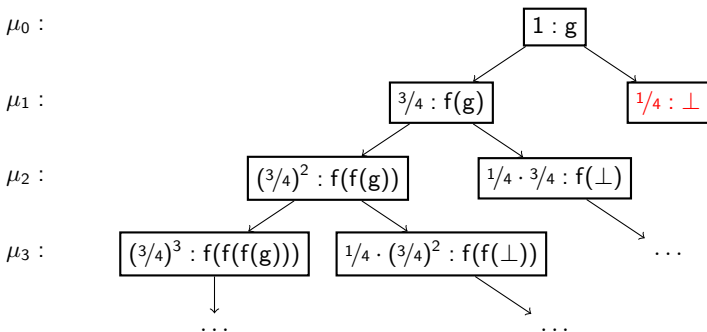


$$\lim_{k \rightarrow \infty} |\mu_k| =$$

Properties for Equality of Termination cont.

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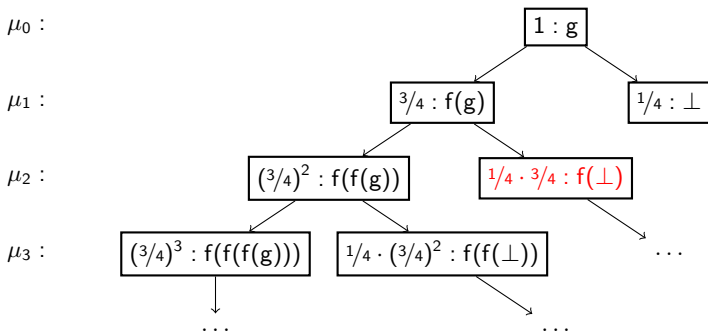
$$\lim_{k \rightarrow \infty} |\mu_k| = 1/4$$

Properties for Equality of Termination cont.

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$$g \rightarrow \{3/4 : f(g), 1/4 : \perp\}$$

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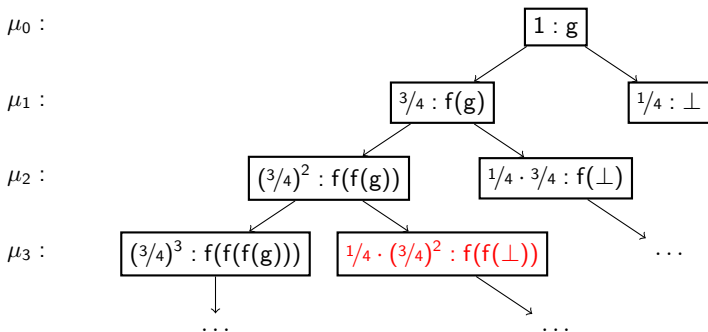
$$\lim_{k \rightarrow \infty} |\mu_k| = 1/4 + 1/4 \cdot 3/4$$

Properties for Equality of Termination cont.

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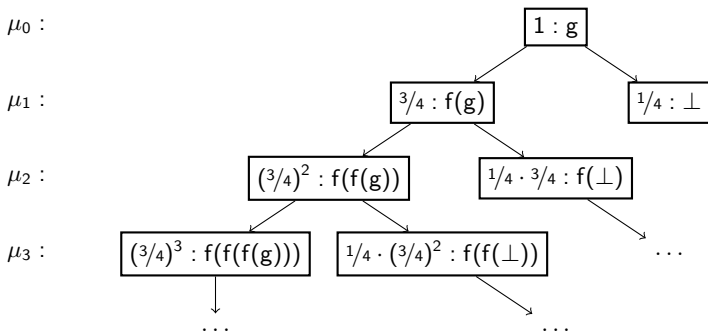


$$\lim_{k \rightarrow \infty} |\mu_k| = 1/4 + 1/4 \cdot 3/4 + 1/4 \cdot (3/4)^2$$

Properties for Equality of Termination cont.

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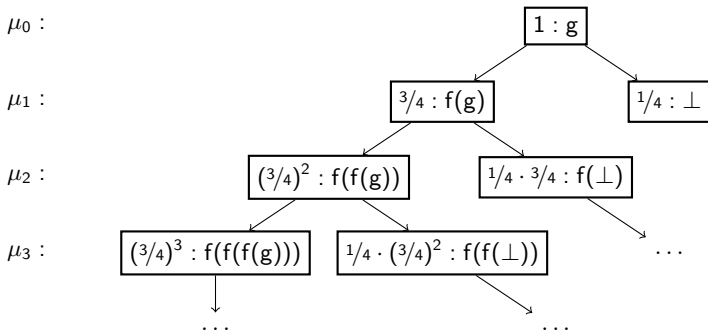


$$\lim_{k \rightarrow \infty} |\mu_k| = 1/4 + 1/4 \cdot 3/4 + 1/4 \cdot (3/4)^2 + \dots$$

Properties for Equality of Termination cont.

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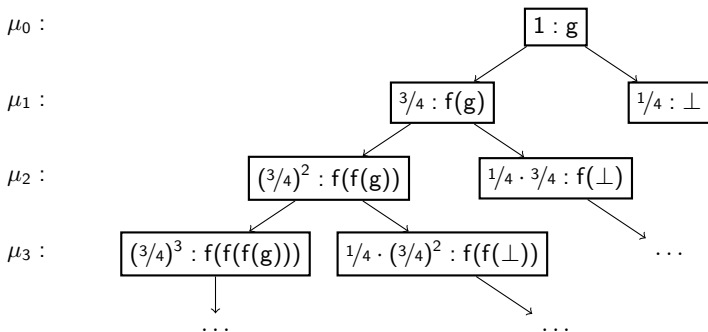


$$\lim_{k \rightarrow \infty} |\mu_k| = 1/4 + 1/4 \cdot 3/4 + 1/4 \cdot (3/4)^2 + \dots = \sum_{i=0}^{\infty} 1/4 \cdot (3/4)^i$$

Properties for Equality of Termination cont.

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$$\lim_{k \rightarrow \infty} |\mu_k| = 1/4 + 1/4 \cdot 3/4 + 1/4 \cdot (3/4)^2 + \dots = \sum_{i=0}^{\infty} 1/4 \cdot (3/4)^i = 1$$

Duplicating

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$\rightarrow \mathcal{S}_2$ is duplicating.

Properties for Equality of AST

Theorem

If \mathcal{S} is non-overlapping and non-duplicating then:

\mathcal{S} is innermost AST iff \mathcal{S} is AST.

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1. Same minimality criterion
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Properties for Equality of AST

Theorem

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1. Same minimality criterion



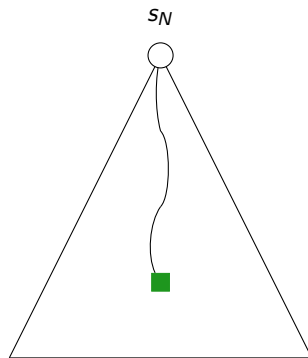
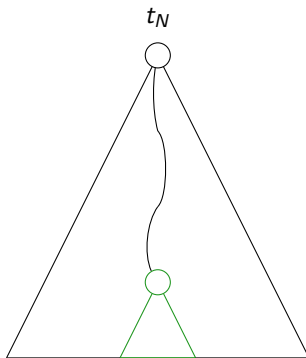
2. Same construction

3. Calculating the probability of termination

Proof cont.

Rewriting possibilities:

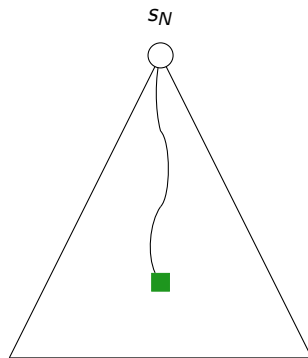
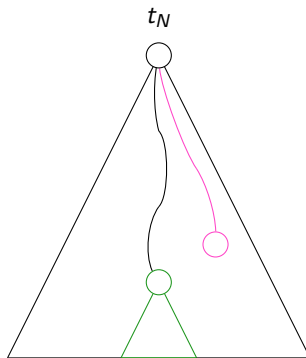
1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position



Proof cont.

Rewriting possibilities:

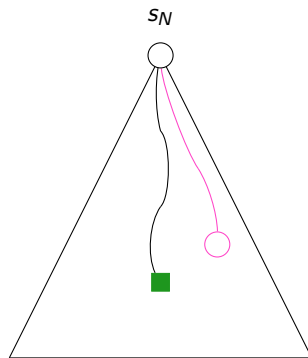
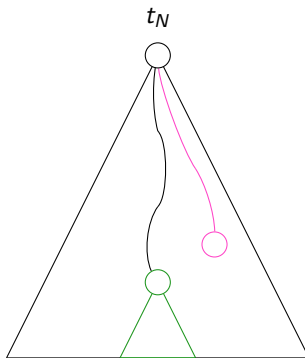
1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position



Proof cont.

Rewriting possibilities:

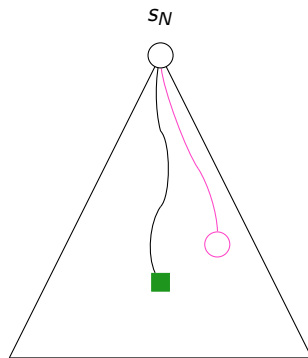
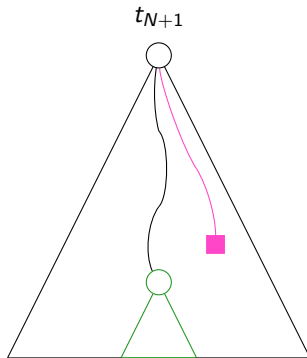
1. Rewriting at an orthogonal position
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Proof cont.

Rewriting possibilities:

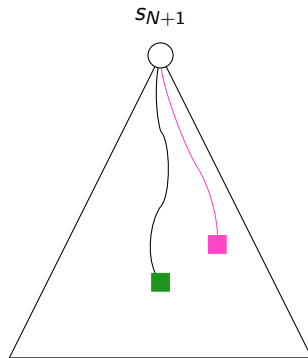
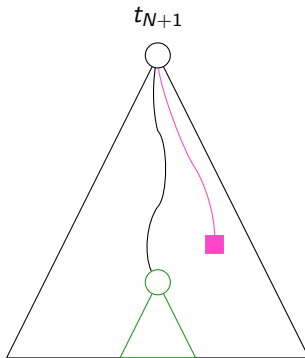
1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position



Proof cont.

Rewriting possibilities:

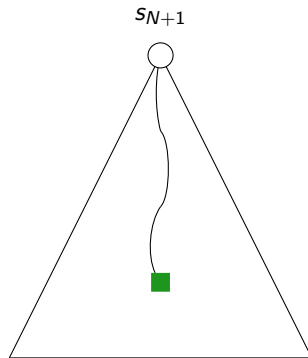
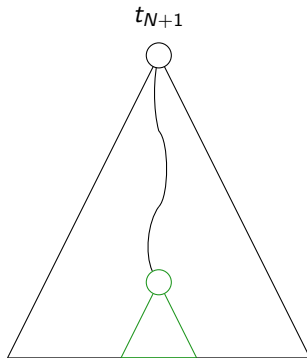
1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position



Proof cont.

Rewriting possibilities:

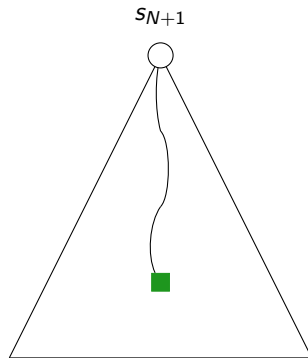
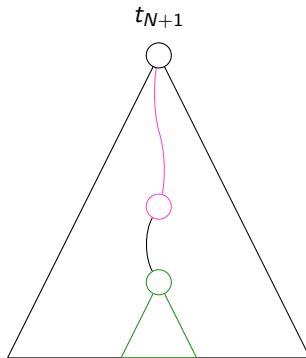
1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position



Proof cont.

Rewriting possibilities:

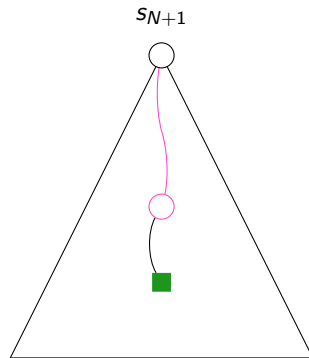
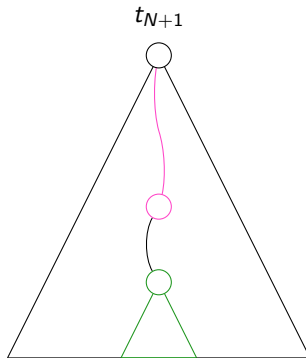
1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position



Proof cont.

Rewriting possibilities:

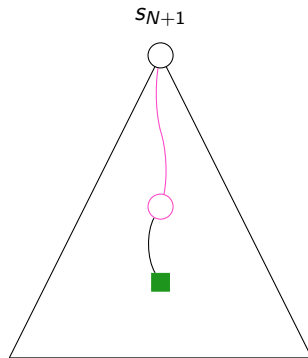
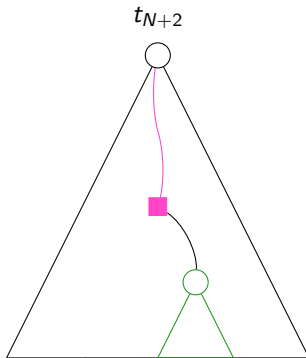
1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position



Proof cont.

Rewriting possibilities:

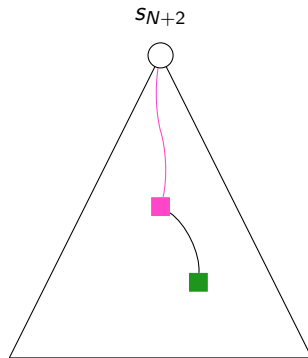
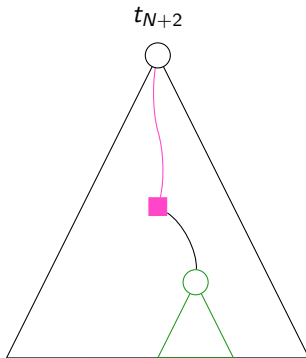
1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position



Proof cont.

Rewriting possibilities:

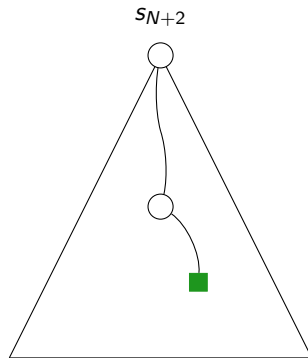
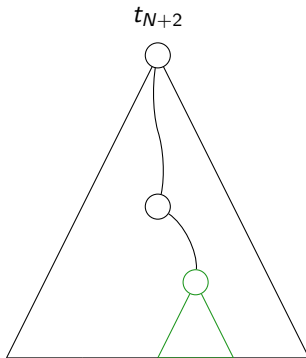
1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position



Proof cont.

Rewriting possibilities:

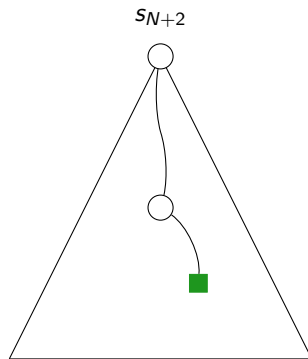
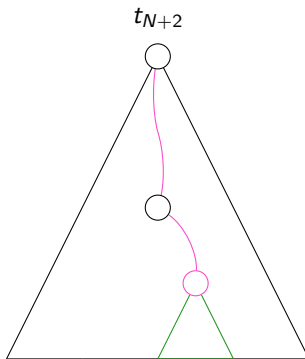
1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position



Proof cont.

Rewriting possibilities:

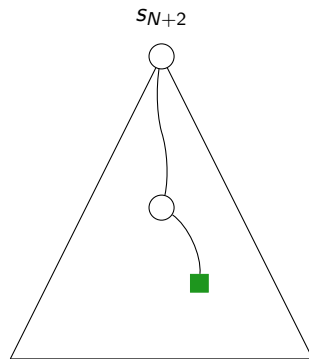
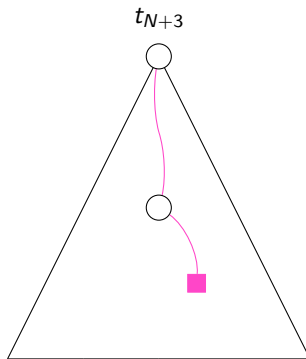
1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position ← only happens once



Proof cont.

Rewriting possibilities:

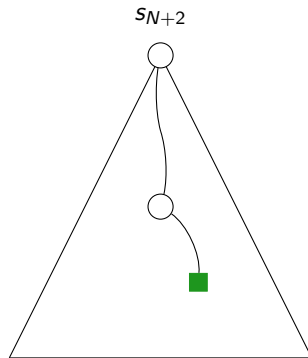
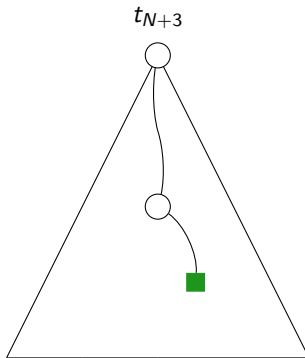
1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position ← only happens once



Proof cont.

Rewriting possibilities:

1. Rewriting at an orthogonal position
2. Rewriting at a position above
3. Rewriting the colored position ← only happens once



Summary

- Structural properties for the equality of AST and innermost AST

Theorem

If \mathcal{S} is non-overlapping and non-duplicating then:

\mathcal{S} is innermost AST iff \mathcal{S} is AST.

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