

Proving Almost-Sure Innermost Termination of Probabilistic Term Rewriting Using Dependency Pairs

Jan-Christoph Kassing, Jürgen Giesl

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Automatic Termination Analysis for TRSs

\mathcal{R}_{plus} :

$$\begin{aligned} \text{plus}(\mathcal{O}, y) &\rightarrow y \\ \text{plus}(s(x), y) &\rightarrow s(\text{plus}(x, y)) \end{aligned}$$

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$$\mathcal{R}_{plus}: \quad \begin{array}{l} plus(\mathcal{O}, y) \rightarrow y \\ plus(s(x), y) \rightarrow s(plus(x, y)) \end{array}$$

Computation “2 + 2”:

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$$\text{plus}(s(s(\mathcal{O})), s(s(\mathcal{O})))$$

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\succ well-founded

There exists no infinite sequence $t_0 \succ t_1 \succ t_2 \succ \dots$

Termination and Complexity Analysis for Programs

Termination and Complexity Analysis for Programs

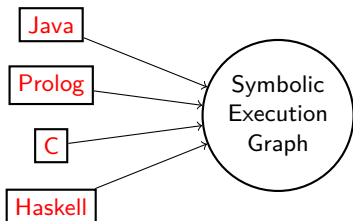
Java

Prolog

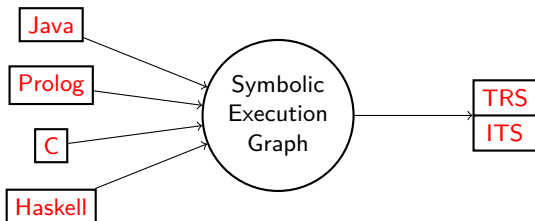
C

Haskell

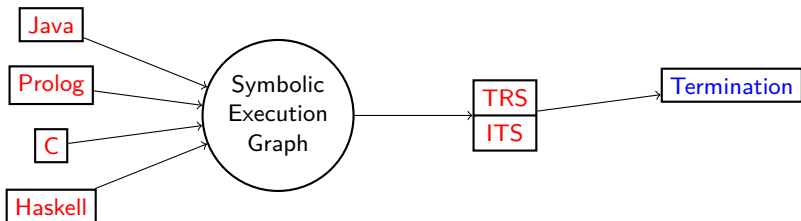
Termination and Complexity Analysis for Programs



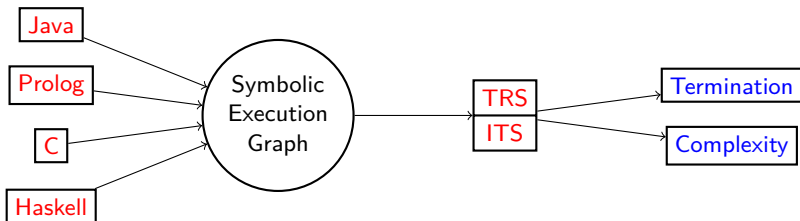
Termination and Complexity Analysis for Programs



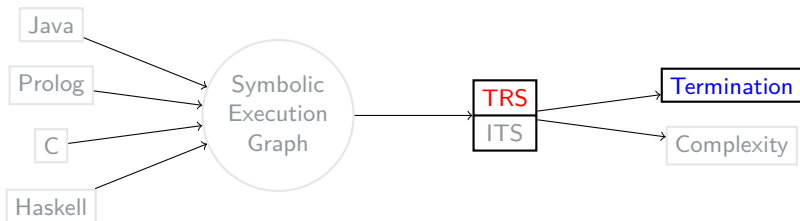
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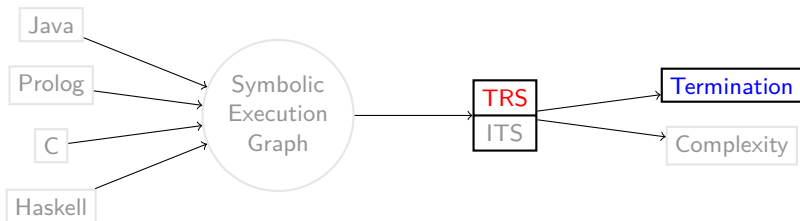


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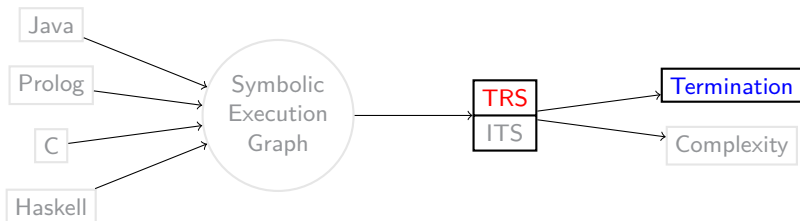
- TRS: especially good for the analysis of algorithms concerning algebraic structures and user-defined data structures

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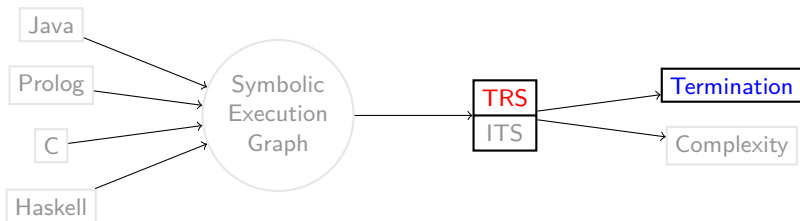


- TRS: especially good for the analysis of algorithms concerning algebraic structures and user-defined data structures
- Turing-complete programming language
⇒ Termination is undecidable

Termination and Complexity Analysis for Programs

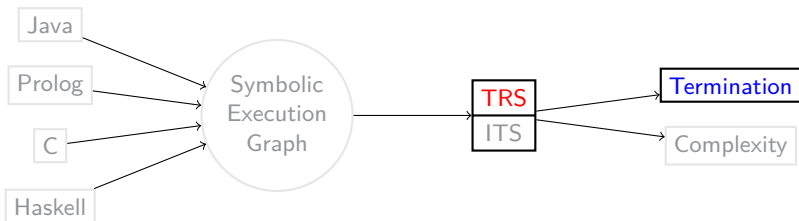


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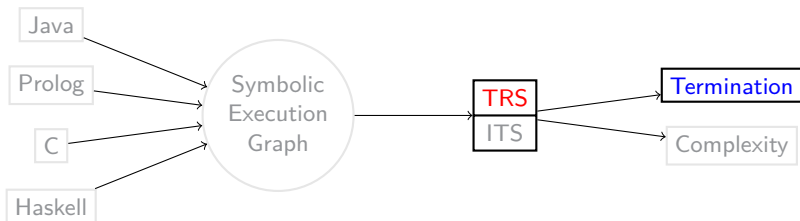
- 1 Direct application of polynomials for termination of TRSs

Termination and Complexity Analysis for Programs



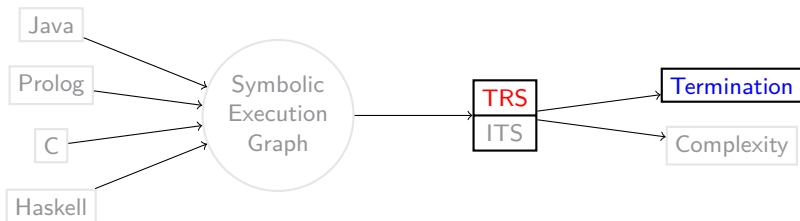
- 1 Direct application of polynomials for termination of TRSs
- 2 DP framework for innermost termination of TRSs

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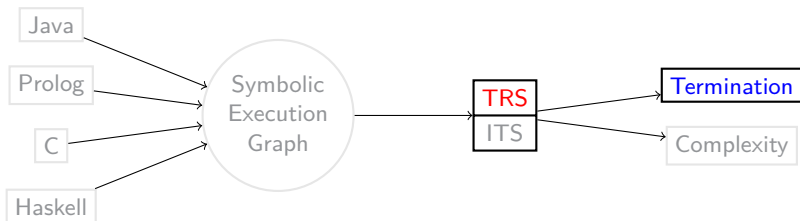
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Termination and Complexity Analysis for Programs



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Termination and Complexity Analysis for Programs



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Automatic Termination Analysis for TRSs [Lankford, 1979]

$$\mathcal{R}_{plus}: \quad \begin{array}{l} \text{plus}(\mathcal{O}, y) \rightarrow y \\ \text{plus}(s(x), y) \rightarrow s(\text{plus}(x, y)) \end{array}$$

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Goal: Find monotonic, natural polynomial interpretation Pol such that

$$\ell \rightarrow r \in \mathcal{R} \text{ implies } Pol(\ell) > Pol(r)$$

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Pol monotonic, natural polynomial interpretation

- natural: $Pol(f) = f_{Pol}$ a polynomial with coefficients $\in \mathbb{N}$
- monotonic: if $x > y$, then $f_{Pol}(\dots, x, \dots) > f_{Pol}(\dots, y, \dots)$

Automatic Termination Analysis for TRSs [Lankford, 1979]

$$\mathcal{R}_{plus}: \quad \begin{array}{l} Pol(\text{plus}(\mathcal{O}, y)) > Pol(y) \\ Pol(\text{plus}(s(x), y)) > Pol(s(\text{plus}(x, y))) \end{array}$$

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$$\begin{aligned} \mathcal{O}_{Pol} &= 0 \\ s_{Pol}(x) &= x + 1 \\ plus_{Pol}(x, y) &= 2x + y + 1 \end{aligned}$$

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Automatic Termination Analysis for TRSs [Lankford, 1979]

$$\mathcal{R}_{plus}: \quad \begin{array}{l} plus_{Pol}(\mathcal{O}_{Pol}, y) > y \\ Pol(plus(s(x), y)) > Pol(s(plus(x, y))) \end{array}$$

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Automatic Termination Analysis for TRSs [Lankford, 1979]

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Automatic Termination Analysis for TRSs [Lankford, 1979]

$$\mathcal{R}_{plus}: \quad \begin{array}{l} y + 1 > y \\ 2x + y + 3 > 2x + y + 2 \end{array}$$

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\Rightarrow proves termination

Dependency Pairs [Arts & Giesl 2000, ...]

\mathcal{R}_{div} :

$minus(x, \mathcal{O})$	\rightarrow	x
$minus(s(x), s(y))$	\rightarrow	$minus(x, y)$
$div(\mathcal{O}, s(y))$	\rightarrow	\mathcal{O}
$div(s(x), s(y))$	\rightarrow	$s(div(minus(x, y), s(y)))$

Dependency Pairs [Arts & Giesl 2000, ...]

$$\begin{array}{l} \mathcal{R}_{div}: \\ \text{minus}(x, \mathcal{O}) \rightarrow x \\ \text{minus}(s(x), s(y)) \rightarrow \text{minus}(x, y) \\ \text{div}(\mathcal{O}, s(y)) \rightarrow \mathcal{O} \\ \text{div}(s(x), s(y)) \rightarrow s(\text{div}(\text{minus}(x, y), s(y))) \end{array}$$

- There exists no monotonic, natural *Pol* that orders all rules strictly

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- There exists no monotonic, natural *PoI* that orders all rules strictly
- Dependency pair approach is able to prove termination

Dependency Pairs [Arts & Giesl 2000, ...]

$$\begin{array}{l} \mathcal{R}_{div}: \\ \quad \text{minus}(x, \mathcal{O}) \rightarrow x \\ \quad \text{minus}(s(x), s(y)) \rightarrow \text{minus}(x, y) \\ \quad \text{div}(\mathcal{O}, s(y)) \rightarrow \mathcal{O} \\ \quad \text{div}(s(x), s(y)) \rightarrow s(\text{div}(\text{minus}(x, y), s(y))) \end{array}$$

Defined Symbols: `minus` and `div`

Dependency Pairs [Arts & Giesl 2000, ...]

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Defined Symbols: `minus` and `div` , **Constructor Symbols:** `s` and `ℒ`

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$\text{Sub}_D(r)$

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root symbol}\}$

Dependency Pairs [Arts & Giesl 2000, ...]

$$\mathcal{R}_{div}: \begin{array}{ll} \text{minus}(x, \mathcal{O}) & \rightarrow x \\ \text{minus}(s(x), s(y)) & \rightarrow \text{minus}(x, y) \\ \text{div}(\mathcal{O}, s(y)) & \rightarrow \mathcal{O} \\ \text{div}(s(x), s(y)) & \rightarrow s(\text{div}(\text{minus}(x, y), s(y))) \end{array}$$

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Sub_D(r)

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root symbol}\}$

$$\begin{array}{ll} \text{Sub}_D(x) & = \emptyset \\ \text{Sub}_D(\text{minus}(x, y)) & = \{\text{minus}(x, y)\} \\ \text{Sub}_D(\mathcal{O}) & = \emptyset \\ \text{Sub}_D(s(\text{div}(\text{minus}(x, y), s(y)))) & = \{\text{minus}(x, y), \text{div}(\text{minus}(x, y), s(y))\} \end{array}$$

Dependency Pairs [Arts & Giesl 2000, ...]

$$\mathcal{R}_{div}: \begin{array}{ll} \text{minus}(x, \mathcal{O}) & \rightarrow x \\ \text{minus}(s(x), s(y)) & \rightarrow \text{minus}(x, y) \\ \text{div}(\mathcal{O}, s(y)) & \rightarrow \mathcal{O} \\ \text{div}(s(x), s(y)) & \rightarrow s(\text{div}(\text{minus}(x, y), s(y))) \end{array}$$

Defined Symbols: `minus` and `div` , **Constructor Symbols:** `s` and `ℒ`

Sub_D(r)

Sub_D(r) := {t | t is a subterm of r with defined root symbol}

Dependency Pairs

If $f(\ell_1, \dots, \ell_n) \rightarrow r$ is a rule and $g(r_1, \dots, r_m) \in \text{Sub}_D(r)$, then $f^\#(\ell_1, \dots, \ell_n) \rightarrow g^\#(r_1, \dots, r_m)$ is a dependency pair

Dependency Pairs [Arts & Giesl 2000, ...]

$$\begin{aligned}
 \mathcal{R}_{div}: \quad & \text{minus}(x, \mathcal{O}) \rightarrow x \\
 & \text{minus}(s(x), s(y)) \rightarrow \text{minus}(x, y) \\
 & \text{div}(\mathcal{O}, s(y)) \rightarrow \mathcal{O} \\
 & \text{div}(s(x), s(y)) \rightarrow s(\text{div}(\text{minus}(x, y), s(y)))
 \end{aligned}$$

Defined Symbols: `minus` and `div` , **Constructor Symbols:** `s` and `ℒ`

$$\begin{aligned}
 \text{Sub}_D(x) &= \emptyset \\
 \text{Sub}_D(\text{minus}(x, y)) &= \{\text{minus}(x, y)\} \\
 \text{Sub}_D(\mathcal{O}) &= \emptyset \\
 \text{Sub}_D(s(\text{div}(\text{minus}(x, y), s(y)))) &= \{\text{minus}(x, y), \text{div}(\text{minus}(x, y), s(y))\}
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$\mathcal{DP}(\mathcal{R}_{div})$:

Dependency Pairs [Arts & Giesl 2000, ...]

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 \end{aligned}$$

$$\begin{aligned}
 \mathcal{DP}(\mathcal{R}_{div}): \\
 M(s(x), s(y)) &\rightarrow M(x, y)
 \end{aligned}$$

Dependency Pairs [Arts & Giesl 2000, ...]

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 \mathcal{R}_{div}: \quad & \text{minus}(x, \mathcal{O}) \rightarrow x \\
 & \text{minus}(s(x), s(y)) \rightarrow \text{minus}(x, y) \\
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 & \text{div}(s(x), s(y)) \rightarrow s(\text{div}(\text{minus}(x, y), s(y)))
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 \mathcal{DP}(\mathcal{R}_{div}): \quad & M(s(x), s(y)) \rightarrow M(x, y) \\
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 \end{aligned}$$

Dependency Pairs [Arts & Giesl 2000, ...]

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 & D(s(x), s(y)) \rightarrow D(\text{minus}(x, y), s(y))
 \end{aligned}$$

Dependency Pairs Cont.

$$\begin{aligned} m(x, \mathcal{O}) &\rightarrow x \\ m(s(x), s(y)) &\rightarrow m(x, y) \\ d(\mathcal{O}, s(y)) &\rightarrow \mathcal{O} \\ d(s(x), s(y)) &\rightarrow s(d(m(x, y), s(y))) \end{aligned}$$

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$(\mathcal{D}, \mathcal{R})$ -Chain

\mathcal{D} a set of DPs, \mathcal{R} a TRS.

A sequence of terms t_0, t_1, \dots is a $(\mathcal{D}, \mathcal{R})$ -chain if we have

$$t_0 \xrightarrow{i}_{\mathcal{D}} \circ \xrightarrow{i}_{\mathcal{R}}^* t_1 \xrightarrow{i}_{\mathcal{D}} \circ \xrightarrow{i}_{\mathcal{R}}^* \dots$$

Dependency Pairs Cont.

$$\begin{aligned} m(x, \mathcal{O}) &\rightarrow x \\ m(s(x), s(y)) &\rightarrow m(x, y) \\ d(\mathcal{O}, s(y)) &\rightarrow \mathcal{O} \\ d(s(x), s(y)) &\rightarrow s(d(m(x, y), s(y))) \end{aligned}$$

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$$D(s^4(\mathcal{O}), s^2(\mathcal{O}))$$

$(DP(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Chain:

Dependency Pairs Cont.

$$\begin{aligned} m(x, \mathcal{O}) &\rightarrow x \\ m(s(x), s(y)) &\rightarrow m(x, y) \\ d(\mathcal{O}, s(y)) &\rightarrow \mathcal{O} \\ d(s(x), s(y)) &\rightarrow s(d(m(x, y), s(y))) \end{aligned}$$

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$(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Chain:

$$\begin{aligned} & \xrightarrow{i}_{\mathcal{DP}(\mathcal{R}_{div})} \\ & D(s^4(\mathcal{O}), s^2(\mathcal{O})) \\ & D(m(s^3(\mathcal{O}), s(\mathcal{O})), s^2(\mathcal{O})) \end{aligned}$$

Dependency Pairs Cont.

$$\begin{aligned} m(x, \mathcal{O}) &\rightarrow x \\ m(s(x), s(y)) &\rightarrow m(x, y) \\ d(\mathcal{O}, s(y)) &\rightarrow \mathcal{O} \\ d(s(x), s(y)) &\rightarrow s(d(m(x, y), s(y))) \end{aligned}$$

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$$\begin{aligned} &\xrightarrow{i}_{\mathcal{DP}(\mathcal{R}_{div})} \\ &\xrightarrow{i^*}_{\mathcal{R}_{div}} \end{aligned} \begin{aligned} &D(s^4(\mathcal{O}), s^2(\mathcal{O})) \\ &D(m(s^3(\mathcal{O}), s(\mathcal{O})), s^2(\mathcal{O})) \\ &D(s^2(\mathcal{O}), s^2(\mathcal{O})) \end{aligned}$$

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$(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Chain:

$$\begin{array}{l} \xrightarrow{i}_{\mathcal{DP}(\mathcal{R}_{div})} \\ \xrightarrow{i}_{\mathcal{R}_{div}}^* \\ \xrightarrow{i}_{\mathcal{DP}(\mathcal{R}_{div})} \end{array} \begin{array}{l} D(s^4(\mathcal{O}), s^2(\mathcal{O})) \\ D(m(s^3(\mathcal{O}), s(\mathcal{O})), s^2(\mathcal{O})) \\ D(s^2(\mathcal{O}), s^2(\mathcal{O})) \\ M(s(\mathcal{O}), s(\mathcal{O})) \end{array}$$

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Dependency Pairs Cont.

$$\begin{aligned} m(x, \mathcal{O}) &\rightarrow x \\ m(s(x), s(y)) &\rightarrow m(x, y) \\ d(\mathcal{O}, s(y)) &\rightarrow \mathcal{O} \\ d(s(x), s(y)) &\rightarrow s(d(m(x, y), s(y))) \end{aligned}$$

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Theorem: Chain Criterion [Arts & Giesl 2000]

\mathcal{R} is innermost terminating iff $(\mathcal{DP}(\mathcal{R}), \mathcal{R})$ is innermost terminating

Dependency Pair Framework

- Key Idea:
 - Transform a “big” problem into simpler sub-problems

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 - DP processors: $Proc(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D}_1, \mathcal{R}_1), \dots, (\mathcal{D}_k, \mathcal{R}_k)\}$

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 - *Proc* is sound: if all $(\mathcal{D}_i, \mathcal{R}_i)$ are innermost terminating, then $(\mathcal{D}, \mathcal{R})$ is innermost terminating
 - *Proc* is complete: if $(\mathcal{D}, \mathcal{R})$ is innermost terminating, then all $(\mathcal{D}_i, \mathcal{R}_i)$ are innermost terminating

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- Processors that reduce \mathcal{D} :

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- Many more...

Dependency Graph Processor (sound & complete)

$$(a) \quad m(x, \mathcal{O}) \rightarrow x$$

$$(b) \quad m(s(x), s(y)) \rightarrow m(x, y)$$

$$(c) \quad d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$$

$$(d) \quad d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$$

$$(1) \quad M(s(x), s(y)) \rightarrow M(x, y)$$

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$$Proc_{DG}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D}_1, \mathcal{R}), \dots, (\mathcal{D}_k, \mathcal{R})\}$$

where $\mathcal{D}_1, \dots, \mathcal{D}_k$ are the SCCs of the $(\mathcal{D}, \mathcal{R})$ -dependency graph:

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$(\mathcal{D}, \mathcal{R})$ -Dependency Graph

- directed graph whose nodes are the dependency pairs from \mathcal{D}

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where $\mathcal{D}_1, \dots, \mathcal{D}_k$ are the SCCs of the $(\mathcal{D}, \mathcal{R})$ -dependency graph:

$(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Dependency Graph:

$$D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$$

$$D(s(x), s(y)) \rightarrow M(x, y)$$

$$M(s(x), s(y)) \rightarrow M(x, y)$$

$(\mathcal{D}, \mathcal{R})$ -Dependency Graph

- directed graph whose nodes are the dependency pairs from \mathcal{D}

Dependency Graph Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

- (1) $M(s(x), s(y)) \rightarrow M(x, y)$
 (2) $D(s(x), s(y)) \rightarrow M(x, y)$
 (3) $D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$

$$\text{Proc}_{DG}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D}_1, \mathcal{R}), \dots, (\mathcal{D}_k, \mathcal{R})\}$$

$$\text{Proc}_{DG}(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$$

where $\mathcal{D}_1, \dots, \mathcal{D}_k$ are the SCCs of the $(\mathcal{D}, \mathcal{R})$ -dependency graph:

$(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Dependency Graph:

$$D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$$

$$D(s(x), s(y)) \rightarrow M(x, y)$$

$$M(s(x), s(y)) \rightarrow M(x, y)$$

$(\mathcal{D}, \mathcal{R})$ -Dependency Graph

- directed graph whose nodes are the dependency pairs from \mathcal{D}
- there is an arc from $s \rightarrow t$ to $v \rightarrow w$ iff $t\sigma_1 \xrightarrow{i}_{\mathcal{R}}^* v\sigma_2$ for substitutions σ_1, σ_2 .

Dependency Graph Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

- (1) $M(s(x), s(y)) \rightarrow M(x, y)$
 (2) $D(s(x), s(y)) \rightarrow M(x, y)$
 (3) $D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$

$$\text{Proc}_{DG}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D}_1, \mathcal{R}), \dots, (\mathcal{D}_k, \mathcal{R})\}$$

$$\text{Proc}_{DG}(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$$

where $\mathcal{D}_1, \dots, \mathcal{D}_k$ are the SCCs of the $(\mathcal{D}, \mathcal{R})$ -dependency graph:

$(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Dependency Graph:

$$D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$$

$$D(s(x), s(y)) \rightarrow M(x, y)$$

$$M(s(x), s(y)) \rightarrow M(x, y)$$

$(\mathcal{D}, \mathcal{R})$ -Dependency Graph

- directed graph whose nodes are the dependency pairs from \mathcal{D}
- there is an arc from $s \rightarrow t$ to $v \rightarrow w$ iff $t\sigma_1 \xrightarrow{i}_{\mathcal{R}}^* v\sigma_2$ for substitutions σ_1, σ_2 .

Dependency Graph Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

- (1) $M(s(x), s(y)) \rightarrow M(x, y)$
 (2) $D(s(x), s(y)) \rightarrow M(x, y)$
 (3) $D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$

$$\text{Proc}_{DG}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D}_1, \mathcal{R}), \dots, (\mathcal{D}_k, \mathcal{R})\}$$

$$\text{Proc}_{DG}(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$$

where $\mathcal{D}_1, \dots, \mathcal{D}_k$ are the SCCs of the $(\mathcal{D}, \mathcal{R})$ -dependency graph:

$(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Dependency Graph:

$$D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$$

$$D(s(x), s(y)) \rightarrow M(x, y)$$

$$M(s(x), s(y)) \rightarrow M(x, y)$$



$(\mathcal{D}, \mathcal{R})$ -Dependency Graph

- directed graph whose nodes are the dependency pairs from \mathcal{D}
- there is an arc from $s \rightarrow t$ to $v \rightarrow w$ iff $t\sigma_1 \xrightarrow{i}^*_{\mathcal{R}} v\sigma_2$ for substitutions σ_1, σ_2 .

Dependency Graph Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

- (1) $M(s(x), s(y)) \rightarrow M(x, y)$
 (2) $D(s(x), s(y)) \rightarrow M(x, y)$
 (3) $D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$

$$\text{Proc}_{DG}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D}_1, \mathcal{R}), \dots, (\mathcal{D}_k, \mathcal{R})\}$$

$$\text{Proc}_{DG}(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$$

where $\mathcal{D}_1, \dots, \mathcal{D}_k$ are the SCCs of the $(\mathcal{D}, \mathcal{R})$ -dependency graph:

$(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Dependency Graph:

$$D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$$

$$D(s(x), s(y)) \rightarrow M(x, y)$$

$$M(s(x), s(y)) \rightarrow M(x, y)$$



$(\mathcal{D}, \mathcal{R})$ -Dependency Graph

- directed graph whose nodes are the dependency pairs from \mathcal{D}
- there is an arc from $s \rightarrow t$ to $v \rightarrow w$ iff $t\sigma_1 \xrightarrow{i}^* v\sigma_2$ for substitutions σ_1, σ_2 .

Dependency Graph Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

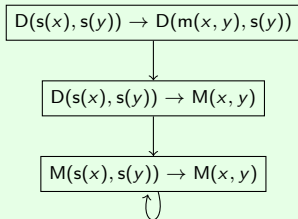
- (1) $M(s(x), s(y)) \rightarrow M(x, y)$
 (2) $D(s(x), s(y)) \rightarrow M(x, y)$
 (3) $D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$

$$\text{Proc}_{DG}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D}_1, \mathcal{R}), \dots, (\mathcal{D}_k, \mathcal{R})\}$$

$$\text{Proc}_{DG}(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$$

where $\mathcal{D}_1, \dots, \mathcal{D}_k$ are the SCCs of the $(\mathcal{D}, \mathcal{R})$ -dependency graph:

$(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Dependency Graph:



$(\mathcal{D}, \mathcal{R})$ -Dependency Graph

- directed graph whose nodes are the dependency pairs from \mathcal{D}
- there is an arc from $s \rightarrow t$ to $v \rightarrow w$ iff $t\sigma_1 \xrightarrow{i}^* v\sigma_2$ for substitutions σ_1, σ_2 .

Dependency Graph Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

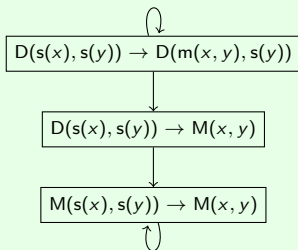
- (1) $M(s(x), s(y)) \rightarrow M(x, y)$
 (2) $D(s(x), s(y)) \rightarrow M(x, y)$
 (3) $D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$

$$\text{Proc}_{DG}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D}_1, \mathcal{R}), \dots, (\mathcal{D}_k, \mathcal{R})\}$$

$$\text{Proc}_{DG}(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$$

where $\mathcal{D}_1, \dots, \mathcal{D}_k$ are the SCCs of the $(\mathcal{D}, \mathcal{R})$ -dependency graph:

$(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Dependency Graph:



$(\mathcal{D}, \mathcal{R})$ -Dependency Graph

- directed graph whose nodes are the dependency pairs from \mathcal{D}
- there is an arc from $s \rightarrow t$ to $v \rightarrow w$ iff $t\sigma_1 \xrightarrow{i}^*_{\mathcal{R}} v\sigma_2$ for substitutions σ_1, σ_2 .

Dependency Graph Processor (sound & complete)

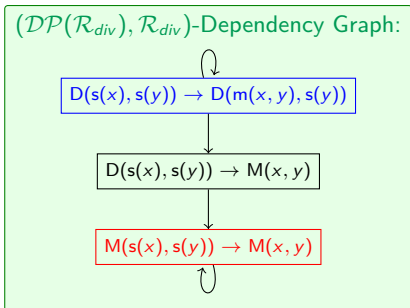
- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

- (1) $M(s(x), s(y)) \rightarrow M(x, y)$
 (2) $D(s(x), s(y)) \rightarrow M(x, y)$
 (3) $D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$

$$Proc_{DG}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D}_1, \mathcal{R}), \dots, (\mathcal{D}_k, \mathcal{R})\}$$

$$Proc_{DG}(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div}) \\ = \{(\{1\}, \mathcal{R}_{div}), (\{3\}, \mathcal{R}_{div})\}$$

where $\mathcal{D}_1, \dots, \mathcal{D}_k$ are the SCCs of the $(\mathcal{D}, \mathcal{R})$ -dependency graph:



$(\mathcal{D}, \mathcal{R})$ -Dependency Graph

- directed graph whose nodes are the dependency pairs from \mathcal{D}
- there is an arc from $s \rightarrow t$ to $v \rightarrow w$ iff $t\sigma_1 \xrightarrow{i}_{\mathcal{R}}^* v\sigma_2$ for substitutions σ_1, σ_2 .

Reduction Pair Processor (sound & complete)

$$(a) \quad m(x, \mathcal{O}) \rightarrow x$$

$$(b) \quad m(s(x), s(y)) \rightarrow m(x, y)$$

$$(c) \quad d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$$

$$(d) \quad d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$$

$$(1) \quad M(s(x), s(y)) \rightarrow M(x, y)$$

$$(2) \quad D(s(x), s(y)) \rightarrow M(x, y)$$

$$(3) \quad D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$$

Reduction Pair Processor (sound & complete)

$$(a) \quad m(x, \mathcal{O}) \rightarrow x$$

$$(b) \quad m(s(x), s(y)) \rightarrow m(x, y)$$

$$(c) \quad d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$$

$$(d) \quad d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$$

$$(1) \quad M(s(x), s(y)) \rightarrow M(x, y)$$

$$(2) \quad D(s(x), s(y)) \rightarrow M(x, y)$$

$$(3) \quad D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$$

Find weakly-monotonic, natural polynomial interpretation Pol

weakly-monotonic

- weakly-monotonic: if $x \geq y$, then $f_{Pol}(\dots, x, \dots) \geq f_{Pol}(\dots, y, \dots)$

Reduction Pair Processor (sound & complete)

$$(a) \quad m(x, \mathcal{O}) \rightarrow x$$

$$(b) \quad m(s(x), s(y)) \rightarrow m(x, y)$$

$$(c) \quad d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$$

$$(d) \quad d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$$

$$(1) \quad M(s(x), s(y)) \rightarrow M(x, y)$$

$$(2) \quad D(s(x), s(y)) \rightarrow M(x, y)$$

$$(3) \quad D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$$

Find weakly-monotonic, natural polynomial interpretation Pol such that

- $Pol(\ell) \geq Pol(r)$ for all rules $\ell \rightarrow r$ in \mathcal{R}
- $Pol(s) > Pol(t)$ for all rules $s \rightarrow t$ in $\mathcal{D}_{>}$
- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Reduction Pair Processor (sound & complete)

$$(a) \quad m(x, \mathcal{O}) \rightarrow x$$

$$(b) \quad m(s(x), s(y)) \rightarrow m(x, y)$$

$$(c) \quad d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$$

$$(d) \quad d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$$

$$(1) \quad M(s(x), s(y)) \rightarrow M(x, y)$$

$$(2) \quad D(s(x), s(y)) \rightarrow M(x, y)$$

$$(3) \quad D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$$

$$Proc_{CRP}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D} \setminus \mathcal{D}_\succ, \mathcal{R})\}$$

Find weakly-monotonic, natural polynomial interpretation Pol such that

- $Pol(\ell) \geq Pol(r)$ for all rules $\ell \rightarrow r$ in \mathcal{R}
- $Pol(s) > Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}_\succ
- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Reduction Pair Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

- (1) $M(s(x), s(y)) \rightarrow M(x, y)$
 (2) $D(s(x), s(y)) \rightarrow M(x, y)$
 (3) $D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$

$$Proc_{RP}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D} \setminus \mathcal{D}_\succ, \mathcal{R})\}$$

$$Proc_{RP}(\{(1)\}, \mathcal{R}_{div})$$

$$Proc_{RP}(\{(3)\}, \mathcal{R}_{div})$$

Find weakly-monotonic, natural polynomial interpretation Pol such that

- $Pol(\ell) \geq Pol(r)$ for all rules $\ell \rightarrow r$ in \mathcal{R}
- $Pol(s) > Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}_\succ
- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Reduction Pair Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

- (1) $M(s(x), s(y)) \rightarrow M(x, y)$
 (2) $D(s(x), s(y)) \rightarrow M(x, y)$
 (3) $D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$

$$Proc_{RP}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D} \setminus \mathcal{D}_{\succ}, \mathcal{R})\}$$

$$Proc_{RP}(\{(1)\}, \mathcal{R}_{div})$$

$$Proc_{RP}(\{(3)\}, \mathcal{R}_{div})$$

$$\begin{aligned} \mathcal{O}_{Pol} &= 0 \\ s_{Pol}(x) &= x + 1 \\ m_{Pol}(x, y) &= x \\ d_{Pol}(x, y) &= x \end{aligned}$$

$(\{(1)\}, \mathcal{R}_{div}) :$

Find weakly-monotonic, natural polynomial interpretation Pol such that

- $Pol(\ell) \geq Pol(r)$ for all rules $\ell \rightarrow r$ in \mathcal{R}
- $Pol(s) > Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}_{\succ}
- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Reduction Pair Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

$$(1) M(s(x), s(y)) \rightarrow M(x, y)$$

$$Proc_{RP}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D} \setminus \mathcal{D}_{\succ}, \mathcal{R})\}$$

$$Proc_{RP}(\{(1)\}, \mathcal{R}_{div})$$

$$Proc_{RP}(\{(3)\}, \mathcal{R}_{div})$$

$$\begin{aligned} \mathcal{O}_{Pol} &= 0 \\ s_{Pol}(x) &= x + 1 \\ m_{Pol}(x, y) &= x \\ d_{Pol}(x, y) &= x \end{aligned}$$

$$(\{(1)\}, \mathcal{R}_{div}) :$$

Find weakly-monotonic, natural polynomial interpretation Pol such that

- $Pol(\ell) \geq Pol(r)$ for all rules $\ell \rightarrow r$ in \mathcal{R}
- $Pol(s) > Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}_{\succ}
- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Reduction Pair Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

$$(1) M(s(x), s(y)) \rightarrow M(x, y)$$

$$Proc_{RP}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D} \setminus \mathcal{D}_\succ, \mathcal{R})\}$$

$$Proc_{RP}(\{(1)\}, \mathcal{R}_{div})$$

$$Proc_{RP}(\{(3)\}, \mathcal{R}_{div})$$

$$\begin{aligned} \mathcal{O}_{Pol} &= 0 \\ s_{Pol}(x) &= x + 1 \\ m_{Pol}(x, y) &= x \\ d_{Pol}(x, y) &= x \end{aligned}$$

$(\{(1)\}, \mathcal{R}_{div}) :$

$$M_{Pol}(x, y) = x$$

Find weakly-monotonic, natural polynomial interpretation Pol such that

- $Pol(\ell) \geq Pol(r)$ for all rules $\ell \rightarrow r$ in \mathcal{R}
- $Pol(s) > Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}_\succ
- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Reduction Pair Processor (sound & complete)

- (a) $Pol(m(x, \mathcal{O})) \geq Pol(x)$
 (b) $Pol(m(s(x), s(y))) \geq Pol(m(x, y))$
 (c) $Pol(d(\mathcal{O}, s(y))) \geq \mathcal{O}$
 (d) $Pol(d(s(x), s(y))) \geq Pol(s(d(m(x, y), s(y))))$

$$(1) Pol(M(s(x), s(y))) > Pol(M(x, y))$$

$$Proc_{RP}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D} \setminus \mathcal{D}_{\succ}, \mathcal{R})\}$$

$$Proc_{RP}(\{(1)\}, \mathcal{R}_{div})$$

$$Proc_{RP}(\{(3)\}, \mathcal{R}_{div})$$

$$\begin{aligned} \mathcal{O}_{Pol} &= 0 \\ s_{Pol}(x) &= x + 1 \\ m_{Pol}(x, y) &= x \\ d_{Pol}(x, y) &= x \end{aligned}$$

$(\{(1)\}, \mathcal{R}_{div}) :$

$$M_{Pol}(x, y) = x$$

Find weakly-monotonic, natural polynomial interpretation Pol such that

- $Pol(\ell) \geq Pol(r)$ for all rules $\ell \rightarrow r$ in \mathcal{R}
- $Pol(s) > Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}_{\succ}
- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Reduction Pair Processor (sound & complete)

$$\begin{aligned} (a) \quad & x \geq x \\ (b) \quad & x + 1 \geq x \\ (c) \quad & 0 \geq 0 \\ (d) \quad & x + 1 \geq x + 1 \end{aligned}$$

$$(1) \quad x + 1 > x$$

$$Proc_{RP}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D} \setminus \mathcal{D}_{>}, \mathcal{R})\}$$

$$Proc_{RP}(\{(1)\}, \mathcal{R}_{div})$$

$$Proc_{RP}(\{(3)\}, \mathcal{R}_{div})$$

$$\begin{aligned} \mathcal{O}_{Pol} &= 0 \\ s_{Pol}(x) &= x + 1 \\ m_{Pol}(x, y) &= x \\ d_{Pol}(x, y) &= x \end{aligned}$$

$$(\{(1)\}, \mathcal{R}_{div}) :$$

$$M_{Pol}(x, y) = x$$

Find weakly-monotonic, natural polynomial interpretation Pol such that

- $Pol(\ell) \geq Pol(r)$ for all rules $\ell \rightarrow r$ in \mathcal{R}
- $Pol(s) > Pol(t)$ for all rules $s \rightarrow t$ in $\mathcal{D}_{>}$
- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Reduction Pair Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

- (1) $M(s(x), s(y)) \rightarrow M(x, y)$
 (2) $D(s(x), s(y)) \rightarrow M(x, y)$
 (3) $D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$

$$Proc_{RP}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D} \setminus \mathcal{D}_{\succ}, \mathcal{R})\}$$

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$$M_{Pol}(x, y) = x$$

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Find weakly-monotonic, natural polynomial interpretation Pol such that

- $Pol(\ell) \geq Pol(r)$ for all rules $\ell \rightarrow r$ in \mathcal{R}
- $Pol(s) > Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}_{\succ}
- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Reduction Pair Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

$$(3) D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$$

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$$(\{(1)\}, \mathcal{R}_{div}) :$$

$$M_{Pol}(x, y) = x$$

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Find weakly-monotonic, natural polynomial interpretation Pol such that

- $Pol(\ell) \geq Pol(r)$ for all rules $\ell \rightarrow r$ in \mathcal{R}
- $Pol(s) > Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}_\succ
- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Reduction Pair Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

$$(3) D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$$

$$Proc_{RP}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D} \setminus \mathcal{D}_\succ, \mathcal{R})\}$$

$$Proc_{RP}(\{(1)\}, \mathcal{R}_{div})$$

$$Proc_{RP}(\{(3)\}, \mathcal{R}_{div})$$

$$\begin{aligned} \mathcal{O}_{Pol} &= 0 \\ s_{Pol}(x) &= x + 1 \\ m_{Pol}(x, y) &= x \\ d_{Pol}(x, y) &= x \end{aligned}$$

$$(\{(1)\}, \mathcal{R}_{div}) :$$

$$M_{Pol}(x, y) = x$$

$$(\{(3)\}, \mathcal{R}_{div}) :$$

$$D_{Pol}(x, y) = x$$

Find weakly-monotonic, natural polynomial interpretation Pol such that

- $Pol(\ell) \geq Pol(r)$ for all rules $\ell \rightarrow r$ in \mathcal{R}
- $Pol(s) > Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}_\succ
- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Reduction Pair Processor (sound & complete)

- (a) $Pol(m(x, \mathcal{O})) \geq Pol(x)$
 (b) $Pol(m(s(x), s(y))) \geq Pol(m(x, y))$
 (c) $Pol(d(\mathcal{O}, s(y))) \geq \mathcal{O}$
 (d) $Pol(d(s(x), s(y))) \geq Pol(s(d(m(x, y), s(y))))$

$$(3) Pol(D(s(x), s(y))) > Pol(D(m(x, y), s(y)))$$

$$Proc_{RP}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D} \setminus \mathcal{D}_{\succ}, \mathcal{R})\}$$

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- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Reduction Pair Processor (sound & complete)

$$\begin{aligned} (a) \quad & x \geq x \\ (b) \quad & x + 1 \geq x \\ (c) \quad & 0 \geq 0 \\ (d) \quad & x + 1 \geq x + 1 \end{aligned}$$

$$Proc_{RP}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D} \setminus \mathcal{D}_{>}, \mathcal{R})\}$$

$$Proc_{RP}(\{(1)\}, \mathcal{R}_{div})$$

$$Proc_{RP}(\{(3)\}, \mathcal{R}_{div})$$

$$(3) \quad x + 1 > x$$

$$\begin{aligned} \mathcal{O}_{Pol} &= 0 \\ s_{Pol}(x) &= x + 1 \\ m_{Pol}(x, y) &= x \\ d_{Pol}(x, y) &= x \end{aligned}$$

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Reduction Pair Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow x$
 (b) $m(s(x), s(y)) \rightarrow m(x, y)$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \mathcal{O}$
 (d) $d(s(x), s(y)) \rightarrow s(d(m(x, y), s(y)))$

- (1) $M(s(x), s(y)) \rightarrow M(x, y)$
 (2) $D(s(x), s(y)) \rightarrow M(x, y)$
 (3) $D(s(x), s(y)) \rightarrow D(m(x, y), s(y))$

$$Proc_{RP}(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D} \setminus \mathcal{D}_{\succ}, \mathcal{R})\}$$

$$Proc_{RP}(\{(1)\}, \mathcal{R}_{div}) = \{(\emptyset, \mathcal{R}_{div})\}$$

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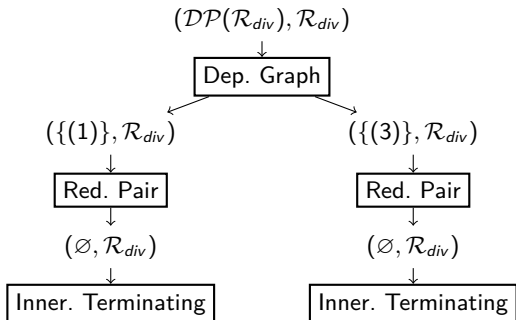
$(\{(3)\}, \mathcal{R}_{div}) :$

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- $Pol(s) \geq Pol(t)$ for all rules $s \rightarrow t$ in \mathcal{D}

Final Innermost Termination Proof



⇒ Innermost termination is proved automatically!

Termination of Probabilistic TRSs

$$\mathcal{R}_{rw}: \quad g(\mathcal{O}) \rightarrow \{1/2 : \mathcal{O}, 1/2 : g(g(\mathcal{O}))\}$$

Termination of Probabilistic TRSs

$$\mathcal{R}_{rw}: \quad g(\mathcal{O}) \rightarrow \{1/2 : \mathcal{O}, 1/2 : g(g(\mathcal{O}))\}$$

Distribution: $\{p_1 : t_1, \dots, p_k : t_k\}$ with $p_1 + \dots + p_k = 1$

Termination of Probabilistic TRSs

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 $\{1 : g(\mathcal{O})\}$

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$$\Rightarrow_{\mathcal{R}_{rw}} \{1/2 : \mathcal{O}, 1/8 : \mathcal{O}, 1/8 : g^2(\mathcal{O}),$$

Termination of Probabilistic TRSs

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$$\Rightarrow_{\mathcal{R}_{rw}} \{1/2 : \mathcal{O}, 1/8 : \mathcal{O}, 1/8 : g^2(\mathcal{O}), 1/8 : g^2(\mathcal{O}), 1/8 : g^4(\mathcal{O})\}$$

Termination of Probabilistic TRSs

$$\mathcal{R}_{rw}: \quad g(\mathcal{O}) \rightarrow \{1/2 : \mathcal{O}, 1/2 : g(g(\mathcal{O}))\}$$

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Termination for PTRSs

[Bournez & Garnier 2005, Avanzini & Dal Lago & Yamada 2019, ...]

- \mathcal{R} is *terminating* iff there is no infinite evaluation $\mu_0 \Rightarrow_{\mathcal{R}} \mu_1 \Rightarrow_{\mathcal{R}} \dots$

Termination of Probabilistic TRSs

$$\mathcal{R}_{rw}: \quad g(\mathcal{O}) \rightarrow \{1/2 : \mathcal{O}, 1/2 : g(g(\mathcal{O}))\}$$

Distribution: $\{p_1 : t_1, \dots, p_k : t_k\}$ with $p_1 + \dots + p_k = 1$
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Termination of Probabilistic TRSs

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Distribution: $\{p_1 : t_1, \dots, p_k : t_k\}$ with $p_1 + \dots + p_k = 1$
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- \mathcal{R} is **almost-surely terminating (AST)**
 iff $\lim_{n \rightarrow \infty} |\mu_n| = 1$ for every infinite evaluation $\mu_0 \Rightarrow_{\mathcal{R}} \mu_1 \Rightarrow_{\mathcal{R}} \dots$

Termination of Probabilistic TRSs

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Distribution: $\{p_1 : t_1, \dots, p_k : t_k\}$ with $p_1 + \dots + p_k = 1$ | μ |

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 $\{1 : g(\mathcal{O})\}$ 0

$$\Rightarrow_{\mathcal{R}_{rw}} \{1/2 : \mathcal{O}, 1/2 : g^2(\mathcal{O})\}$$

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Termination of Probabilistic TRSs

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$$\{1 : g(\mathcal{O})\} \quad 0$$

$$\Rightarrow_{\mathcal{R}_{rw}} \{1/2 : \mathcal{O}, 1/2 : g^2(\mathcal{O})\} \quad 1/2$$

$$\Rightarrow_{\mathcal{R}_{rw}} \{1/2 : \mathcal{O}, 1/4 : g(\mathcal{O}), 1/4 : g^3(\mathcal{O})\}$$

$$\Rightarrow_{\mathcal{R}_{rw}} \{1/2 : \mathcal{O}, 1/8 : \mathcal{O}, 1/8 : g^2(\mathcal{O}), 1/8 : g^2(\mathcal{O}), 1/8 : g^4(\mathcal{O})\}$$

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$\{1 : g(\mathcal{O})\}$ 0

$\Rightarrow_{\mathcal{R}_{rw}} \{1/2 : \mathcal{O}, 1/2 : g^2(\mathcal{O})\}$ 1/2

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Termination for PTRSs

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$\{1 : g(\mathcal{O})\}$ 0

$\Rightarrow_{\mathcal{R}_{rw}} \{1/2 : \mathcal{O}, 1/2 : g^2(\mathcal{O})\}$ 1/2

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Termination of Probabilistic TRSs

$$\mathcal{R}_{rw} : \quad g(x) \rightarrow \left\{ \frac{1}{2} : x, \frac{1}{2} : g(g(x)) \right\}$$

Theorem (AST with Polynomial Interpretation)

Let Pol be a **multilinear monotonic polynomial interpretation**.

For all $\ell \rightarrow \mu = \{p_1 : r_1, \dots, p_k : r_k\} \in \mathcal{R}$ let

- $Pol(\ell) > Pol(r_j)$ for some $1 \leq j \leq k$
- $Pol(\ell) \geq \mathbb{E}(\mu) = p_1 \cdot Pol(r_1) + \dots + p_k \cdot Pol(r_k)$

Then \mathcal{R} is AST.

Termination of Probabilistic TRSs

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Pol is multilinear

monomials like $x \cdot y$, but no monomials like x^2

Termination of Probabilistic TRSs

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$$g_{Pol}(x) = 1 + x$$

Termination of Probabilistic TRSs

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Termination of Probabilistic TRSs

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$$g_{Pol}(x) = 1 + x$$

Termination of Probabilistic TRSs

$$\mathcal{R}_{rw} : 1 + x \geq \frac{1}{2} \cdot x + \frac{1}{2} \cdot (2 + x)$$

Theorem (AST with Polynomial Interpretation)

Let Pol be a **multilinear monotonic polynomial interpretation**.

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Termination of Probabilistic TRSs

$$\mathcal{R}_{rw} : 1 + x \geq 1 + x$$

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Termination of Probabilistic TRSs

$$\mathcal{R}_{rw} : 1 + x \geq 1 + x$$

Theorem (AST with Polynomial Interpretation)

Let Pol be a **multilinear monotonic polynomial interpretation**.

For all $\ell \rightarrow \mu = \{p_1 : r_1, \dots, p_k : r_k\} \in \mathcal{R}$ let

- $Pol(\ell) > Pol(r_j)$ for some $1 \leq j \leq k$
- $Pol(\ell) \geq \mathbb{E}(\mu) = p_1 \cdot Pol(r_1) + \dots + p_k \cdot Pol(r_k)$

Then \mathcal{R} is AST.

Pol is multilinear

monomials like $x \cdot y$, but no monomials like x^2

$$g_{Pol}(x) = 1 + x$$

\Rightarrow proves AST

Dependency Pairs for AST: Failed Attempt

$\text{Sub}_D(r)$

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root}\}$

Dependency Pairs for AST: Failed Attempt

$\text{Sub}_D(r)$

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root}\}$

Options for Dependency Pairs (A)

If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, then a dependency pair is :

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Options for Dependency Pairs (A)

If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, then a dependency pair is :

(A) : $\{ \ell^\# \rightarrow \{p_1 : t_1^\#, \dots, p_j : t_j^\#, \dots, p_k : t_k^\#\} \mid t_j \in \text{Sub}_D(r_j), 1 \leq i \leq k \}$

Dependency Pairs for AST: Failed Attempt

$\text{Sub}_D(r)$

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root}\}$

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If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, then a dependency pair is :

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If $\text{Sub}_D(r_i) = \emptyset$, then we insert a constructor \perp into $\text{Sub}_D(r_i)$

Dependency Pairs for AST: Failed Attempt

Sub_D(r)

Sub_D(r) := {t | t is a subterm of r with defined root}

Options for Dependency Pairs (A)

If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, then a dependency pair is :

(A) : { $\ell^\# \rightarrow \{p_1 : t_1^\#, \dots, p_j : t_j^\#, \dots, p_k : t_k^\#\} \mid t_j \in \text{Sub}_D(r_j), 1 \leq i \leq k$ }

If $\text{Sub}_D(r_i) = \emptyset$, then we insert a constructor \perp into $\text{Sub}_D(r_i)$

\mathcal{R}_1 : g $\rightarrow \{1/2 : f(g, g), 1/2 : \perp\}$ AST

\mathcal{R}_2 : g $\rightarrow \{1/2 : f(g, g, g), 1/2 : \perp\}$ not AST

Dependency Pairs for AST: Failed Attempt

Sub_D(r)

Sub_D(r) := {t | t is a subterm of r with defined root}

Options for Dependency Pairs (A)

If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, then a dependency pair is :

$$(A) : \{ \ell^\# \rightarrow \{p_1 : t_1^\#, \dots, p_j : t_j^\#, \dots, p_k : t_k^\#\} \mid t_j \in \text{Sub}_D(r_j), 1 \leq i \leq k \}$$

If $\text{Sub}_D(r_i) = \emptyset$, then we insert a constructor \perp into $\text{Sub}_D(r_i)$

$$\begin{array}{lll} \mathcal{R}_1 & : g & \rightarrow \{1/2 : f(g, g), 1/2 : \perp\} & \text{AST} \\ \text{DP}(\mathcal{R}_1) & : G & \rightarrow \{1/2 : G, 1/2 : \perp\} & \text{AST} \end{array}$$

$$\mathcal{R}_2 : g \rightarrow \{1/2 : f(g, g, g), 1/2 : \perp\} \quad \text{not AST}$$

Dependency Pairs for AST: Failed Attempt

Sub_D(r)

Sub_D(r) := {t | t is a subterm of r with defined root}

Options for Dependency Pairs (A)

If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, then a dependency pair is :

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If $\text{Sub}_D(r_i) = \emptyset$, then we insert a constructor \perp into $\text{Sub}_D(r_i)$

\mathcal{R}_1	: g	$\rightarrow \{1/2 : f(g, g), 1/2 : \perp\}$	AST
DP(\mathcal{R}_1)	: G	$\rightarrow \{1/2 : G, 1/2 : \perp\}$	AST

\mathcal{R}_2	: g	$\rightarrow \{1/2 : f(g, g, g), 1/2 : \perp\}$	not AST
DP(\mathcal{R}_2)	: G	$\rightarrow \{1/2 : G, 1/2 : \perp\}$	AST ⚡

Dependency Tuples for AST: Failed Attempt

$\text{Sub}_D(r)$

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root}\}$ - as multiset

Dependency Tuples for AST: Failed Attempt

$\text{Sub}_D(r)$

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root}\}$ - as multiset

Options for Dependency Tuples (B)

If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, and $\text{Sub}_D(r_j) = \{t_{j,1}, \dots, t_{j,i_j}\}$ for all $1 \leq j \leq k$, then the dependency pair is:

(B) : $\ell^\# \rightarrow \{p_1 : \text{Com}(t_{1,1}^\#, \dots, t_{1,i_1}^\#), \dots, p_k : \text{Com}(t_{k,1}^\#, \dots, t_{k,i_k}^\#)\}$

Dependency Tuples for AST: Failed Attempt

$\text{Sub}_D(r)$

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root}\}$ - **as multiset**

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If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, and $\text{Sub}_D(r_j) = \{t_{j,1}, \dots, t_{j,i_j}\}$ for all $1 \leq j \leq k$, then the dependency pair is:

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\mathcal{R}_1 : $g \rightarrow \{1/2 : f(g, g), 1/2 : \perp\}$ AST

\mathcal{R}_2 : $g \rightarrow \{1/2 : f(g, g, g), 1/2 : \perp\}$ not AST

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If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, and $\text{Sub}_D(r_j) = \{t_{j,1}, \dots, t_{j,i_j}\}$ for all $1 \leq j \leq k$, then the dependency pair is:

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\mathcal{R}_1	: g	$\rightarrow \{1/2 : f(g, g), 1/2 : \perp\}$	AST
$DP(\mathcal{R}_1)$: G	$\rightarrow \{1/2 : \text{Com}(G, G), 1/2 : \perp\}$	AST
\mathcal{R}_2	: g	$\rightarrow \{1/2 : f(g, g, g), 1/2 : \perp\}$	not AST

Dependency Tuples for AST: Failed Attempt

$\text{Sub}_D(r)$

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root}\}$ - **as multiset**

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If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, and $\text{Sub}_D(r_j) = \{t_{j,1}, \dots, t_{j,i_j}\}$ for all $1 \leq j \leq k$, then the dependency pair is:

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\mathcal{R}_1	: g	$\rightarrow \{1/2 : f(g, g), 1/2 : \perp\}$	AST
$DP(\mathcal{R}_1)$: G	$\rightarrow \{1/2 : \text{Com}(G, G), 1/2 : \perp\}$	AST
\mathcal{R}_2	: g	$\rightarrow \{1/2 : f(g, g, g), 1/2 : \perp\}$	not AST
$DP(\mathcal{R}_2)$: G	$\rightarrow \{1/2 : \text{Com}(G, G, G), 1/2 : \perp\}$	not AST

Dependency Tuples for AST: Failed Attempt

$\text{Sub}_D(r)$

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root}\}$ - **as multiset**

Options for Dependency Tuples (B)

If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, and $\text{Sub}_D(r_j) = \{t_{j,1}, \dots, t_{j,i_j}\}$ for all $1 \leq j \leq k$, then the dependency pair is:

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$\mathcal{R}_3 :$

$f(\mathcal{O})$	\rightarrow	$\{1 : f(\mathbf{a})\},$
\mathbf{a}	\rightarrow	$\{1/2 : \mathbf{b}, 1/2 : \mathbf{c}\}$

Dependency Tuples for AST: Failed Attempt

$\text{Sub}_D(r)$

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root}\}$ - as multiset

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If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, and $\text{Sub}_D(r_j) = \{t_{j,1}, \dots, t_{j,i_j}\}$ for all $1 \leq j \leq k$, then the dependency pair is:

(B) : $\ell^\# \rightarrow \{p_1 : \text{Com}(t_{1,1}^\#, \dots, t_{1,i_1}^\#), \dots, p_k : \text{Com}(t_{k,1}^\#, \dots, t_{k,i_k}^\#)\}$

$$\begin{aligned} \mathcal{R}_3 : \quad & f(\mathcal{O}) \rightarrow \{1 : f(\mathbf{a})\}, \\ & \mathbf{a} \rightarrow \{1/2 : \mathbf{b}, 1/2 : \mathbf{c}\} \\ \text{DT}(\mathcal{R}_3) : \quad & F(\mathcal{O}) \rightarrow \{1 : \text{Com}(F(\mathbf{a}), \mathbf{A})\} \end{aligned}$$

Dependency Tuples for AST: Failed Attempt

$\text{Sub}_D(r)$

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root}\}$ - **as multiset**

Options for Dependency Tuples (B)

If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, and $\text{Sub}_D(r_j) = \{t_{j,1}, \dots, t_{j,i_j}\}$ for all $1 \leq j \leq k$, then the dependency pair is:

(B) : $\ell^\# \rightarrow \{p_1 : \text{Com}(t_{1,1}^\#, \dots, t_{1,i_1}^\#), \dots, p_k : \text{Com}(t_{k,1}^\#, \dots, t_{k,i_k}^\#)\}$

$$\begin{array}{lll} \mathcal{R}_3 : & f(\mathcal{O}) & \rightarrow \{1 : f(a)\}, \\ & a & \rightarrow \{1/2 : b, 1/2 : c\} \\ DT(\mathcal{R}_3) : & F(\mathcal{O}) & \rightarrow \{1 : \text{Com}(F(a), A)\} \\ & A & \rightarrow \{1/2 : B, 1/2 : C\} \end{array}$$

Dependency Tuples for AST: Failed Attempt

$$\begin{array}{l} \mathcal{R}_3 : \quad f(\mathcal{O}) \rightarrow \{1 : f(a)\}, \\ \quad \quad a \rightarrow \{1/2 : b, 1/2 : c\} \\ DT(\mathcal{R}_3) : F(\mathcal{O}) \rightarrow \{1 : \text{Com}(F(a), A)\} \\ \quad \quad A \rightarrow \{1/2 : B, 1/2 : C\} \end{array}$$

$$\{ 1 : f(\mathcal{O}) \}$$

Dependency Tuples for AST: Failed Attempt

$$\begin{array}{lcl}
 \mathcal{R}_3 : & f(\mathcal{O}) & \rightarrow \{1 : f(a)\}, \\
 & a & \rightarrow \{1/2 : b, 1/2 : c\} \\
 DT(\mathcal{R}_3) : & F(\mathcal{O}) & \rightarrow \{1 : \text{Com}(F(a), A)\} \\
 & A & \rightarrow \{1/2 : B, 1/2 : C\}
 \end{array}$$

$$\begin{array}{l}
 \stackrel{i}{\Rightarrow}_{\mathcal{R}_3} \quad \{ 1 : f(\mathcal{O}) \} \\
 \quad \quad \quad \{ 1 : f(a) \}
 \end{array}$$

Dependency Tuples for AST: Failed Attempt

$$\begin{array}{l}
 \mathcal{R}_3 : \quad f(\mathcal{O}) \rightarrow \{1 : f(\mathbf{a})\}, \\
 \quad \quad \mathbf{a} \rightarrow \{1/2 : \mathbf{b}, 1/2 : \mathbf{c}\} \\
 DT(\mathcal{R}_3) : \quad F(\mathcal{O}) \rightarrow \{1 : \text{Com}(F(\mathbf{a}), \mathbf{A})\} \\
 \quad \quad \quad \mathbf{A} \rightarrow \{1/2 : \mathbf{B}, 1/2 : \mathbf{C}\}
 \end{array}$$

$$\begin{array}{l}
 \xrightarrow{i} \\
 \xrightarrow{i} \mathcal{R}_3 \\
 \xrightarrow{i} \\
 \xrightarrow{i} \mathcal{R}_3
 \end{array}
 \quad
 \begin{array}{l}
 \{ 1 : f(\mathcal{O}) \} \\
 \{ 1 : f(\mathbf{a}) \} \\
 \{ 1/2 : f(\mathbf{b}), 1/2 : f(\mathbf{c}) \}
 \end{array}$$

Dependency Tuples for AST: Failed Attempt

$$\begin{array}{l}
 \mathcal{R}_3 : \quad f(\mathcal{O}) \rightarrow \{1 : f(\mathbf{a})\}, \\
 \quad \quad \quad \mathbf{a} \rightarrow \{1/2 : \mathbf{b}, 1/2 : \mathbf{c}\} \\
 DT(\mathcal{R}_3) : \quad F(\mathcal{O}) \rightarrow \{1 : \text{Com}(F(\mathbf{a}), \mathbf{A})\} \\
 \quad \quad \quad \mathbf{A} \rightarrow \{1/2 : \mathbf{B}, 1/2 : \mathbf{C}\}
 \end{array}$$

$$\begin{array}{l}
 \begin{array}{l} \overset{i}{\Rightarrow} \\ \overset{i}{\Rightarrow} \\ \Rightarrow \end{array} \mathcal{R}_3 \quad \begin{array}{l} \{ 1 : f(\mathcal{O}) \} \\ \{ 1 : f(\mathbf{a}) \} \\ \{ 1/2 : f(\mathbf{b}), 1/2 : f(\mathbf{c}) \} \end{array} \\
 \\
 \{ 1 : F(\mathcal{O}) \}
 \end{array}$$

Dependency Tuples for AST: Failed Attempt

$$\begin{array}{l}
 \mathcal{R}_3 : \quad f(\mathcal{O}) \rightarrow \{1 : f(a)\}, \\
 \quad \quad a \rightarrow \{1/2 : b, 1/2 : c\} \\
 DT(\mathcal{R}_3) : \quad F(\mathcal{O}) \rightarrow \{1 : \text{Com}(F(a), A)\} \\
 \quad \quad A \rightarrow \{1/2 : B, 1/2 : C\}
 \end{array}$$

$$\begin{array}{l}
 \{ 1 : f(\mathcal{O}) \} \\
 \xrightarrow{i} \mathcal{R}_3 \quad \{ 1 : f(a) \} \\
 \xrightarrow{i} \mathcal{R}_3 \quad \{ 1/2 : f(b), 1/2 : f(c) \}
 \end{array}$$

$$\begin{array}{l}
 \{ 1 : F(\mathcal{O}) \} \\
 \xrightarrow{i} DT(\mathcal{R}_3) \quad \{ 1 : \text{Com}(F(a), A) \}
 \end{array}$$

Dependency Tuples for AST: Failed Attempt

$$\begin{array}{l}
 \mathcal{R}_3 : \quad f(\mathcal{O}) \rightarrow \{1 : f(a)\}, \\
 \quad \quad a \rightarrow \{1/2 : b, 1/2 : c\} \\
 DT(\mathcal{R}_3) : \quad F(\mathcal{O}) \rightarrow \{1 : \text{Com}(F(a), A)\} \\
 \quad \quad \quad A \rightarrow \{1/2 : B, 1/2 : C\}
 \end{array}$$

$$\begin{array}{l}
 \{ 1 : f(\mathcal{O}) \} \\
 \xrightarrow{i} \mathcal{R}_3 \quad \{ 1 : f(a) \} \\
 \xrightarrow{i} \mathcal{R}_3 \quad \{ 1/2 : f(b), 1/2 : f(c) \}
 \end{array}$$

$$\begin{array}{l}
 \{ 1 : F(\mathcal{O}) \} \\
 \xrightarrow{i} DT(\mathcal{R}_3) \quad \{ 1 : \text{Com}(F(a), A) \} \\
 \xrightarrow{i} DT(\mathcal{R}_3) \quad \{ 1/2 : \text{Com}(F(a), B), 1/2 : \text{Com}(F(a), C) \}
 \end{array}$$

Dependency Tuples for AST: Failed Attempt

$$\begin{array}{l}
 \mathcal{R}_3 : \quad f(\mathcal{O}) \rightarrow \{1 : f(a)\}, \\
 \quad \quad \quad \mathbf{a} \rightarrow \{1/2 : \mathbf{b}, 1/2 : \mathbf{c}\} \\
 DT(\mathcal{R}_3) : \quad F(\mathcal{O}) \rightarrow \{1 : \text{Com}(F(a), A)\} \\
 \quad \quad \quad A \rightarrow \{1/2 : B, 1/2 : C\}
 \end{array}$$

$$\begin{array}{l}
 \{ 1 : f(\mathcal{O}) \} \\
 \xrightarrow{i} \mathcal{R}_3 \quad \{ 1 : f(a) \} \\
 \xrightarrow{i} \mathcal{R}_3 \quad \{ 1/2 : f(\mathbf{b}), 1/2 : f(\mathbf{c}) \}
 \end{array}$$

$$\begin{array}{l}
 \{ 1 : F(\mathcal{O}) \} \\
 \xrightarrow{i} DT(\mathcal{R}_3) \quad \{ 1 : \text{Com}(F(a), A) \} \\
 \xrightarrow{i} DT(\mathcal{R}_3) \quad \{ 1/2 : \text{Com}(F(\mathbf{a}), B), 1/2 : \text{Com}(F(a), C) \} \\
 \xrightarrow{i} \mathcal{R}_3 \quad \{ 1/4 : \text{Com}(F(\mathbf{b}), B), 1/4 : \text{Com}(F(\mathbf{c}), B), \}
 \end{array}$$

Dependency Tuples for AST: Failed Attempt

$$\begin{array}{l}
 \mathcal{R}_3 : \quad f(\mathcal{O}) \rightarrow \{1 : f(a)\}, \\
 \quad \quad \quad a \rightarrow \{1/2 : b, 1/2 : c\} \\
 DT(\mathcal{R}_3) : \quad F(\mathcal{O}) \rightarrow \{1 : \text{Com}(F(a), A)\} \\
 \quad \quad \quad A \rightarrow \{1/2 : B, 1/2 : C\}
 \end{array}$$

$$\begin{array}{l}
 \{ 1 : f(\mathcal{O}) \} \\
 \xrightarrow{i} \mathcal{R}_3 \quad \{ 1 : f(a) \} \\
 \xrightarrow{i} \mathcal{R}_3 \quad \{ 1/2 : f(b), 1/2 : f(c) \}
 \end{array}$$

$$\begin{array}{l}
 \{ 1 : F(\mathcal{O}) \} \\
 \xrightarrow{i} DT(\mathcal{R}_3) \quad \{ 1 : \text{Com}(F(a), A) \} \\
 \xrightarrow{i} DT(\mathcal{R}_3) \quad \{ 1/2 : \text{Com}(F(a), B), 1/2 : \text{Com}(F(a), C) \} \\
 \xrightarrow{i} \mathcal{R}_3 \quad \{ 1/4 : \text{Com}(F(b), B), 1/4 : \text{Com}(F(c), B), \\
 \quad \quad \quad 1/4 : \text{Com}(F(b), C), 1/4 : \text{Com}(F(c), C) \}
 \end{array}$$

Dependency Triples for AST: Failed Attempt

$$\begin{array}{l}
 \mathcal{R}_3 : \quad f(\mathcal{O}) \rightarrow \{1 : f(a)\}, \\
 \quad \quad a \rightarrow \{1/2 : b, 1/2 : c\} \\
 DT(\mathcal{R}_3) : \quad F(\mathcal{O}) \rightarrow \{1 : \text{Com}(F(a), A)\} \\
 \quad \quad A \rightarrow \{1/2 : B, 1/2 : C\}
 \end{array}$$

$$\begin{array}{l}
 \begin{array}{l} i \\ \Rightarrow_{\mathcal{R}_3} \\ i \\ \Rightarrow_{\mathcal{R}_3} \end{array} \quad \{ 1 : f(\mathcal{O}) \} \\
 \{ 1 : f(a) \} \\
 \{ 1/2 : f(b), 1/2 : f(c) \}
 \end{array}$$

$$\begin{array}{l}
 \begin{array}{l} i \\ \Rightarrow_{DT(\mathcal{R}_3)} \\ i \\ \Rightarrow_{DT(\mathcal{R}_3)} \\ i \\ \Rightarrow_{\mathcal{R}_3} \end{array} \quad \{ 1 : F(\mathcal{O}) \} \\
 \{ 1 : \text{Com}(F(a), A) \} \\
 \{ 1/2 : \text{Com}(F(a), B), 1/2 : \text{Com}(F(a), C) \} \\
 \{ 1/4 : \text{Com}(F(b), B), 1/4 : \text{Com}(F(c), B), \\
 \quad 1/4 : \text{Com}(F(b), C), 1/4 : \text{Com}(F(c), C) \}
 \end{array}$$

- The **red** terms do not correspond to a term in the original rewrite sequence
- One cannot simulate original rewrite sequences by chains

Coupled Dependency Tuples: Sound

$\text{Sub}_D(r)$

$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root}\}$ - as multiset

Coupled Dependency Tuples: Sound

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Options for Dependency Tuples (C)

If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, and $\text{Sub}_D(r_j) = \{t_{1,j}, \dots, t_{i_j,j}\}$ for all $1 \leq j \leq k$, then a dependency pair is:

$$(C): \quad \ell^\# \quad \rightarrow \{p_1 : \text{Com}(t_{1,1}^\#, \dots, t_{1,i_1}^\#) \quad , \dots , p_k : \text{Com}(t_{k,1}^\#, \dots, t_{k,i_k}^\#) \quad \}$$

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If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, and $\text{Sub}_D(r_j) = \{t_{1,j}, \dots, t_{i_j,j}\}$ for all $1 \leq j \leq k$, then a dependency pair is:

$$(C): \langle \ell^\#, \ell \rangle \rightarrow \{p_1 : \langle \text{Com}(t_{1,1}^\#, \dots, t_{1,i_1}^\#), r_1 \rangle, \dots, p_k : \langle \text{Com}(t_{k,1}^\#, \dots, t_{k,i_k}^\#), r_k \rangle\}$$

Coupled Dependency Tuples: Sound

$\text{Sub}_D(r)$

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If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, and $\text{Sub}_D(r_j) = \{t_{1,j}, \dots, t_{i_j,j}\}$ for all $1 \leq j \leq k$, then a dependency pair is:

(C): $\langle \ell^\#, \ell \rangle \rightarrow \{p_1 : \langle \text{Com}(t_{1,1}^\#, \dots, t_{1,i_1}^\#), r_1 \rangle, \dots, p_k : \langle \text{Com}(t_{k,1}^\#, \dots, t_{k,i_k}^\#), r_k \rangle\}$

$\mathcal{R}_3 :$

$f(\mathcal{O})$	$\rightarrow \{1 : f(a)\},$
a	$\rightarrow \{1/2 : b, 1/2 : c\}$

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If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, and $\text{Sub}_D(r_j) = \{t_{1,j}, \dots, t_{i_j,j}\}$ for all $1 \leq j \leq k$, then a dependency pair is:

(C): $\langle \ell^\#, \ell \rangle \rightarrow \{p_1 : \langle \text{Com}(t_{1,1}^\#, \dots, t_{1,i_1}^\#), r_1 \rangle, \dots, p_k : \langle \text{Com}(t_{k,1}^\#, \dots, t_{k,i_k}^\#), r_k \rangle\}$

$$\begin{aligned} \mathcal{R}_3 : \quad & f(\mathcal{O}) && \rightarrow \{1 : f(a)\}, \\ & a && \rightarrow \{1/2 : b, 1/2 : c\} \\ DT(\mathcal{R}_3) : \quad & \langle F(\mathcal{O}), f(\mathcal{O}) \rangle && \rightarrow \{1 : \langle \text{Com}(F(a), A), f(a) \rangle\} \end{aligned}$$

Coupled Dependency Tuples: Sound

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$\text{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root}\}$ - as multiset

Options for Dependency Tuples (C)

If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is a rule, and $\text{Sub}_D(r_j) = \{t_{1,j}, \dots, t_{i_j,j}\}$ for all $1 \leq j \leq k$, then a dependency pair is:

(C): $\langle \ell^\#, \ell \rangle \rightarrow \{p_1 : \langle \text{Com}(t_{1,1}^\#, \dots, t_{1,i_1}^\#), r_1 \rangle, \dots, p_k : \langle \text{Com}(t_{k,1}^\#, \dots, t_{k,i_k}^\#), r_k \rangle\}$

$$\begin{array}{lll} \mathcal{R}_3 : & f(\mathcal{O}) & \rightarrow \{1 : f(a)\}, \\ & a & \rightarrow \{1/2 : b, 1/2 : c\} \\ DT(\mathcal{R}_3) : & \langle F(\mathcal{O}), f(\mathcal{O}) \rangle & \rightarrow \{1 : \langle \text{Com}(F(a), A), f(a) \rangle\} \\ & \langle A, a \rangle & \rightarrow \{1/2 : \langle B, b \rangle, 1/2 : \langle C, c \rangle\} \end{array}$$

Coupled Dependency Tuples: Sound

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 & & \{ 1 : f(\mathcal{O}) \}
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 \end{array}$$

$$\begin{array}{l}
 \xrightarrow{i} \\
 \xrightarrow{\mathcal{R}_3}
 \end{array}
 \quad
 \begin{array}{l}
 \{ 1 : f(\mathcal{O}) \} \\
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Coupled Dependency Tuples: Sound

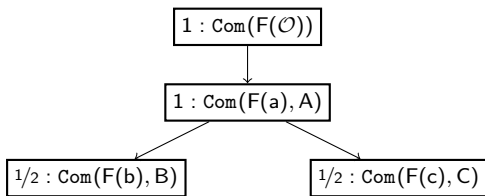
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Probabilistic Chain

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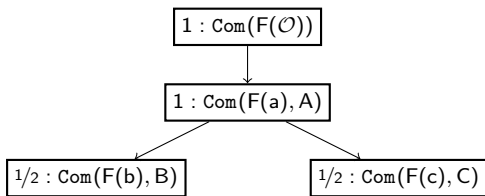


Chain: On every path, we use a dependency pair after a finite number of steps.

$$(\overset{i}{\rightarrow}_D \circ \overset{i}{\rightarrow}_R^*)$$

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Chain: On every path, we use a dependency pair after a finite number of steps.

$$(\rightarrow_D \circ \rightarrow_{\mathcal{R}}^*)$$

Theorem: Chain Criterion

\mathcal{R} is innermost AST if $(DP(\mathcal{R}), \mathcal{R})$ is innermost AST

Dependency Tuples for \mathcal{R}_{div}

\mathcal{R}_{div} :

- (1) $m(x, \mathcal{O}) \rightarrow \{1 : x\}$
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Dependency Pair Framework for Proving iAST of PTRS

- Our objects we work with:
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- Again, many more...

Dependency Graph Processor (sound & complete)

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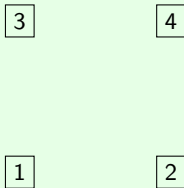
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$(DP(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Dependency Graph:



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- directed graph whose nodes are the dependency tuples from \mathcal{P}

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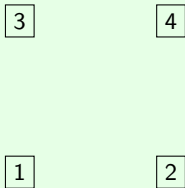
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Dependency Graph Processor (sound & complete)

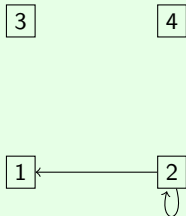
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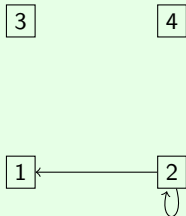
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$(\mathcal{P}, \mathcal{S})$ -Dependency Graph

- directed graph whose nodes are the dependency tuples from \mathcal{P}
- there is an arc from $s \rightarrow \{p_1 : c_1, \dots, p_k : c_k\}$ to $v \rightarrow \dots$ iff there is $t \triangleleft c_j$ for some j and substitutions σ_1, σ_2 such that $t\sigma_1 \xrightarrow{i}_{np(\mathcal{S})}^* v\sigma_2$

Dependency Graph Processor (sound & complete)

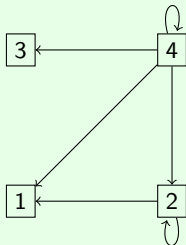
- (a) $m(x, \mathcal{O}) \rightarrow \{1 : x\}$
 (b) $m(s(x), s(y)) \rightarrow \{1 : m(x, y)\}$
 (c) $d(\mathcal{O}, s(y)) \rightarrow \{1 : \mathcal{O}\}$
 (d) $d(s(x), s(y)) \rightarrow \{1/2 : d(s(x), s(y)), 1/2 : s(d(m(x, y), s(y)))\}$

- (1) $M(x, \mathcal{O}) \rightarrow \{1 : \text{Com}\}$
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$$\text{Proc}_{DG}(\mathcal{P}, \mathcal{S}) = \{(\mathcal{P}_1, \mathcal{S}), \dots, (\mathcal{P}_k, \mathcal{S})\}$$

where $\mathcal{P}_1, \dots, \mathcal{P}_k$ are the SCCs of the $(\mathcal{P}, \mathcal{S})$ -dependency graph

$(DP(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Dependency Graph:



$(\mathcal{P}, \mathcal{S})$ -Dependency Graph

- directed graph whose nodes are the dependency tuples from \mathcal{P}
- there is an arc from $s \rightarrow \{p_1 : c_1, \dots, p_k : c_k\}$ to $v \rightarrow \dots$ iff there is $t \triangleleft c_j$ for some j and substitutions σ_1, σ_2 such that $t\sigma_1 \xrightarrow{i}_{np(\mathcal{S})}^* v\sigma_2$

Dependency Graph Processor (sound & complete)

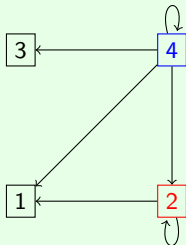
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$$\text{Proc}_{DG}(\mathcal{DT}(\mathcal{R}_{div}), \mathcal{R}_{div})$$

$$= \{(\{(2)\}, \mathcal{R}_{div}), (\{(4)\}, \mathcal{R}_{div})\}$$

$(\mathcal{DP}(\mathcal{R}_{div}), \mathcal{R}_{div})$ -Dependency Graph:



$(\mathcal{P}, \mathcal{S})$ -Dependency Graph

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- there is an arc from $s \rightarrow \{p_1 : c_1, \dots, p_k : c_k\}$ to $v \rightarrow \dots$ iff there is $t \triangleleft c_j$ for some j and substitutions σ_1, σ_2 such that $t\sigma_1 \xrightarrow{i}_{\text{np}(\mathcal{S})}^* v\sigma_2$

Reduction Pair Processor (sound & complete)

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Reduction Pair Processor (sound & complete)

- | | | | |
|-----|--|-----|--|
| (a) | $m(x, \mathcal{O}) \rightarrow \{1 : x\}$ | (1) | $M(x, \mathcal{O}) \rightarrow \{1 : \text{Com}\}$ |
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$$\text{Proc}_{RP}(\mathcal{P}, \mathcal{S}) = \{(\mathcal{P} \setminus \mathcal{P}_{\succ}, \mathcal{S})\}$$

Find weakly-monotonic, multilinear, Com-additive, natural polynomial interpretation *Pol* such that

Reduction Pair Processor (sound & complete)

- | | |
|--|--|
| (a) $m(x, \mathcal{O}) \rightarrow \{1 : x\}$ | (1) $M(x, \mathcal{O}) \rightarrow \{1 : \text{Com}\}$ |
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$$\text{Proc}_{RP}(\mathcal{P}, \mathcal{S}) = \{(\mathcal{P} \setminus \mathcal{P}_\succ, \mathcal{S})\}$$

Find **weakly-monotonic**, **multilinear**, **Com-additive**, **natural polynomial interpretation** *Pol* such that

- For all $\ell \rightarrow \mu = \{p_1 : r_1, \dots, p_k : r_k\}$ in \mathcal{S} :

$$\text{Pol}(\ell) \geq \mathbb{E}(\mu) = \sum_{1 \leq j \leq k} p_j \cdot \text{Pol}(r_j)$$

Reduction Pair Processor (sound & complete)

- | | |
|--|--|
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- For all $\langle \ell^\#, \ell \rangle \rightarrow \mu = \{p_1 : \langle c_1, r_1 \rangle, \dots, p_k : \langle c_k, r_k \rangle\}$ in \mathcal{P} :

$$\text{Pol}(\ell^\#) \geq \mathbb{E}(\mu) = \sum_{1 \leq j \leq k} p_j \cdot \text{Pol}(c_j)$$

Reduction Pair Processor (sound & complete)

- | | |
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| (a) $m(x, \mathcal{O}) \rightarrow \{1 : x\}$ | (1) $M(x, \mathcal{O}) \rightarrow \{1 : \text{Com}\}$ |
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$$\text{Proc}_{RP}(\mathcal{P}, \mathcal{S}) = \{(\mathcal{P} \setminus \mathcal{P}_{>}, \mathcal{S})\}$$

Find **weakly-monotonic, multilinear, Com-additive, natural polynomial interpretation** *Pol* such that

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- For all $\langle \ell^\#, \ell \rangle \rightarrow \{p_1 : \langle c_1, r_1 \rangle, \dots, p_k : \langle c_k, r_k \rangle\}$ in $\mathcal{P}_{>}$ there is a j with

$$\text{Pol}(\ell^\#) > \text{Pol}(c_j)$$

If $\ell \rightarrow \{p_1 : r_1, \dots, p_k : r_k\}$ is in \mathcal{S} , then we additionally require

$$\text{Pol}(\ell) \geq \text{Pol}(r_j)$$

Reduction Pair Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow \{1 : x\}$
 (b) $m(s(x), s(y)) \rightarrow \{1 : m(x, y)\}$
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$(\{(4)\}, \mathcal{R}_{div}) :$

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$(\{(4)\}, \mathcal{R}_{div}) :$

\mathcal{O}_{Pol}	$=$	0	$s_{Pol}(x)$	$=$	$2x + 2$
$m_{Pol}(x, y)$	$=$	x	$d_{Pol}(x, y)$	$=$	x
$M_{Pol}(x, y)$	$=$	$x + 1$	$D_{Pol}(x, y)$	$=$	$x + 1$
$\text{Com}_{Pol}(x, y)$	$=$	$x + y$			

Reduction Pair Processor (sound & complete)

- (a) $m(x, \mathcal{O}) \rightarrow \{1 : x\}$
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 M_{Pol}(x, y) & = x + 1 & D_{Pol}(x, y) & = x + 1 \\
 \text{Com}_{Pol}(x, y) & = x + y & &
 \end{array}$$

$$\begin{aligned}
 Pol(D(s(x), s(y))) & \geq 1/2 \cdot Pol(D(s(x), s(y))) \\
 & \quad + 1/2 \cdot Pol(\text{Com}(D(m(x, y), s(y)), M(x, y)))
 \end{aligned}$$

Reduction Pair Processor (sound & complete)

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 \end{array}$$

$$2x + 3 \geq 1/2 \cdot (2x + 3) + 1/2 \cdot (2x + 2)$$

Reduction Pair Processor (sound & complete)

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 \end{array}$$

$$2x + 3 \geq 2x + 2 + 1/2$$

Reduction Pair Processor (sound & complete)

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$$2x + 3 \geq 2x + 2 + 1/2$$

and

$$Pol(D(s(x), s(y))) = 2x + 3 > 2x + 2 = Pol(\text{Com}(D(m(x, y), s(y)), M(x, y)))$$

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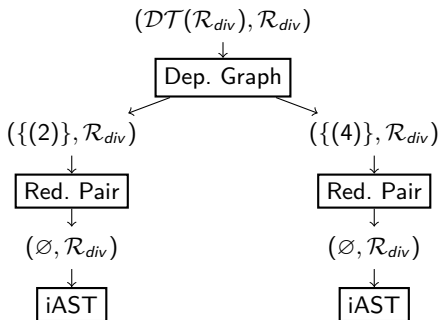
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$$ProcRP(\{(4)\}, \mathcal{R}_{div}) = \{(\emptyset, \mathcal{R}_{div})\}$$

Final Innermost Almost-Sure Termination Proof



⇒ Innermost almost-sure termination is proved automatically!

Implementation and Experiments

- Fully implemented in AProVE
- Evaluated on 67 benchmarks (61 iAST / 59 AST)

	AProVE	DPs	Direct Polo	NaTT2
iAST	53	51	27	22
AST	27	-	27	22

Probabilistic Quicksort:

$$\text{rotate}(\text{cons}(x, xs)) \rightarrow \{1/2 : \text{cons}(x, xs), 1/2 : \text{rotate}(\text{app}(xs, \text{cons}(x, \text{nil})))\}$$

$$\text{qs}(\text{nil}) \rightarrow \{1 : \text{nil}\}$$

$$\text{qs}(\text{cons}(x, xs)) \rightarrow \{1 : \text{qsHelp}(\text{rotate}(\text{cons}(x, xs)))\}$$

$$\text{qsHelp}(\text{cons}(x, xs)) \rightarrow \{1 : \text{app}(\text{qs}(\text{low}(x, xs)), \text{cons}(x, \text{qs}(\text{high}(x, xs))))\}$$

...

Conclusion

1. Direct application of polynomials for AST of probabilistic TRSs

- $Pol(\ell) > Pol(r_j)$ for some $1 \leq j \leq k$
- $Pol(\ell) \geq \mathbb{E}(\mu) = p_1 \cdot Pol(r_1) + \dots + p_k \cdot Pol(r_k)$

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2. DP framework for innermost AST of probabilistic TRSs

- New Dependency Tuples and Chains:

$$\langle \ell^\#, \ell \rangle \rightarrow \{p_1 : \langle \text{Com}(t_{1,1}^\#, \dots, t_{1,i_1}^\#), r_1 \rangle, \dots, p_k : \langle \text{Com}(t_{k,1}^\#, \dots, t_{k,i_k}^\#), r_k \rangle\}$$

Conclusion

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- Adapted the main processors and added more:
 - Dependency Graph Processor
 - Usable Terms Processor
 - Reduction Pair Processor
 - Usable Rules Processor
 - Probability Removal Processor

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